

# Random Dirichlet Polynomials: Metric Entropy Approach

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# 1. Random processes and metric entropy

Random process:

$$X(t), \quad t \in T.$$

Natural distance on  $T$ :

$$\rho(s, t)^2 = \mathbb{E}(X(s) - X(t))^2.$$

Covering numbers and metric entropy

$$N(\varepsilon) = \inf\{n : \exists A_1 \dots A_n : T \subset \cup_1^n A_j, \text{diam} A_j \leq \varepsilon\}.$$

$$H(\varepsilon) = \log N(\varepsilon).$$

Main idea:

Size of  $X \leftrightarrow$  growth of  $H(\varepsilon)$  as  $\varepsilon \rightarrow 0$ .

## 2. Basic results using metric entropy

Let  $X$  be centered:  $\mathbb{E}X(t) = 0$  and Gaussian:  $X(t)$  has a normal (Gaussian) distribution.

We measure the size of  $X$  with  $\mathbb{E} \sup_{t \in T} X(t)$ .  
Then:

a) *Sudakov–Fernique lower bound*

$$\mathbb{E} \sup_{t \in T} X(t) \geq C \varepsilon H(\varepsilon)^{1/2} \quad \forall \varepsilon > 0.$$

( $C = 1.5$ )

b) *Dudley upper bound*

$$\mathbb{E} \sup_{t \in T} X(t) \leq C \int_0^\infty H(\varepsilon)^{1/2} d\varepsilon$$

( $C = 4\sqrt{2}$ )

### 3. Dirichlet series and polynomials

Dirichlet series:

$$\varphi(t) = \sum_{n=1}^{\infty} a_n e^{it\lambda_n}, \quad t \in \mathbb{R}.$$

Dirichlet polynomials:

$$\mathcal{P}_N(t) = \sum_{n=1}^N a_n e^{it\lambda_n}, \quad t \in \mathbb{R}.$$

Trigonometric series and polynomials (Fourier theory):  $\lambda_n = n$ .

Classical Dirichlet series and polynomials:  $\lambda_n = \log n$ .

$$\varphi(t) = \sum_{n=1}^{\infty} a_n n^{it}, \quad t \in \mathbb{R}.$$

Example: Riemann zeta function  $a_n = n^{-\sigma}$ ,

$$\varphi(t) = \sum_{n=1}^{\infty} n^{-\sigma+it}, \quad t \in \mathbb{R}.$$

General task: extend Fourier theory to Dirichlet case.

## A small subtask: Rudin-Shapiro polynomials

Consider a trigonometric polynomial:

$$\mathcal{P}(t) = \sum_{n=0}^N a_n e^{itn}, \quad 0 \leq t \leq 2\pi.$$

Then,

$$\|\mathcal{P}\|_\infty = \sup_t |\mathcal{P}(t)| \leq \sum_n |a_n|.$$

Conversely,

$$\|\mathcal{P}\|_\infty^2 \geq \frac{1}{2\pi} \int_0^{2\pi} |\mathcal{P}(t)|^2 dt = \sum_n |a_n|^2 \geq \left( \frac{\sum_n |a_n|}{\sqrt{N}} \right)^2.$$

Thus

$$\|\mathcal{P}\|_\infty \geq \frac{\sum_n |a_n|}{\sqrt{N}}.$$

Rudin and Shapiro:  $\exists(a_n) \in \{\pm 1\}$ :

$$\|\mathcal{P}\|_\infty \sim (2 + \sqrt{2})\sqrt{N} \approx \frac{\sum_n |a_n|}{\sqrt{N}}.$$

These  $\mathcal{P}$  are *small deviation* polynomials.

What about small deviation Dirichlet polynomials ? ( $-\infty < t < \infty$ ). Konyagin (Moscow) and Queffelec (Lille) proved:

a)  $\exists c_1, \beta_1 : \forall N, \forall (a_n)$

$$\|\mathcal{P}\|_\infty \geq \frac{c_1 \sum_n |a_n|}{\sqrt{N}} \exp\left(\beta_1 \sqrt{\log N \cdot \log \log N}\right) .$$

b)  $\exists c_2, \beta_2, N_0 : \forall N \geq N_0, \exists \mathcal{P}_N = \sum_{n=1}^N a_n n^{it} :$

$$\|\mathcal{P}\|_\infty \leq \frac{c_2 \sum_n |a_n|}{\sqrt{N}} \exp\left(\beta_2 \sqrt{\log N \cdot \log \log N}\right) .$$

This is attained by choosing *random* coefficients and studying  $\mathbb{E}\|\mathcal{P}_N\|_\infty$ .

Rudin-Shapiro bound is not attained.

No explicit  $(a_n)$  is known.

Complicated answer.

## 5. Size of random Dirichlet polynomials

Take the simplest classical random Dirichlet polynomials

$$\mathcal{P}_N(t, \omega) = \sum_{n=1}^N n^{-\sigma} \varepsilon_n(\omega) n^{it}$$

with  $\varepsilon_n$  independent standard normal or symmetric Bernoulli ( $\pm 1$ ).

**Theorem 1** (*Halász, Queffelec*) *For any positive  $\sigma < 1$*

$$M_N := \mathbb{E} \|\mathcal{P}_N(t, \omega)\|_\infty \approx \frac{N^{1-\sigma}}{\log N}, \quad N \rightarrow \infty.$$

Since  $\sum_{n=1}^N |a_n| \approx N^{1-\sigma}$ , we see that  $M_N$  is smaller than  $\sum_{n=1}^N |a_n|$  but much larger than  $\frac{\sum_{n=1}^N |a_n|}{\sqrt{N}}$ . These random polynomials are not the small deviation ones.

We will consider somewhat more tricky polynomials.

Let  $(p_\tau)$  be the sequence of all prime numbers.

Denote  $P^+(n)$  the largest prime divisor of  $n$ .

Denote

$$\mathcal{V}_{\tau,N} = \{n : n \leq N, P^+(n) \leq p_\tau\} \subset [1; N]$$

and consider random polynomial

$$\mathcal{P}_{\tau,N}(t) = \sum_{n \in \mathcal{V}_{\tau,N}} n^{-\sigma} \varepsilon_n n^{it}$$

as a candidate for being a small deviation polynomial. Then

$$M_{\tau,N} = \mathbb{E} \sup_{t \in \mathbb{R}} |\mathcal{P}_{\tau,N}(t)|$$

should be compared with  $\sum_{n \in \mathcal{V}_{\tau,N}} n^{-\sigma}$ .

Important:  $a_n = n^{-\sigma}$  satisfy  $a_{mn} = a_m a_n$ .

## 6. Number theory: Dickman function.

In order to understand  $\sum_{n \in \mathcal{V}_{\tau, N}} n^{-\sigma}$ , let  $\sigma = 0$  and consider  $|\mathcal{V}_{\tau, N}| = |\{n : P^+(n) \leq p_{\tau}, n \leq N\}|$ . This is related to Dickman function.

If  $M \rightarrow \infty$ ,  $N \rightarrow \infty$  so that

$$\frac{\log N}{\log M} \rightarrow u \quad (M \approx N^{1/u}),$$

then

$$\frac{|\{n : P^+(n) \leq M, n \leq N\}|}{N} \rightarrow \rho(u).$$

(there is a differential equation but no explicit formula for  $\rho(u)$ ).

Therefore, if  $\tau \approx N^{1/u}$ , then  $|\mathcal{V}_{\tau, N}| \sim \rho(u)N$ .

## 7. Main result Recall that

$$M_{\tau,N} = \mathbb{E} \sup_{t \in \mathbb{R}} |\mathcal{P}_{\tau,N}(t)|$$

**Theorem 2** *For any  $\sigma < 1$ , the average  $M_{\tau,N}$  satisfies:*

$$M_{\tau,N} \leq \begin{cases} C \frac{N^{1/2-\sigma} \tau^{1/2}}{(\log N)^{1/2}}, & N^{1/2} \leq \tau \leq N, \\ C \frac{N^{3/4-\sigma}}{(\log N)^{1/2}}, & \frac{N^{1/2}}{\log N} \leq \tau \leq N^{1/2}, \\ C N^{1/2-\sigma} \tau^{1/2} \rho\left(\frac{N}{p_\tau}\right)^{1/2}, & \tau \leq \frac{N^{1/2}}{\log N} \end{cases}$$

and

$$M_{\tau,N} \geq C \frac{N^{1/2-\sigma} \tau^{1/2}}{(\log \tau)^{1/2}} \rho\left(\frac{N}{p_\tau}\right)^{1/2}.$$

**Remark:** in the first zone the estimates are order-optimal. In other zones a log-gap appears. First zone contains Halász-Queffelec estimate.

**Open problem:** find the size asymptotics of random Dirichlet polynomials

$$\mathbb{E} \sup_{t \in \mathbb{R}} \left| \sum_{n=1}^N a_n \varepsilon_n n^{it}(t) \right| \approx ???$$

for suitable class of multiplicative  $(a_n)$ .

The upper bound in Theorem 2 with the right choice of  $\tau$ ,

$$\tau \approx \exp \left\{ \sqrt{\frac{\log N}{2} \log \log N} \right\}$$

yields (as by Konyagin and Queffelec)

$$\|\mathcal{P}_{\tau, N}\|_{\infty} \leq \frac{c_2 \sum_n |a_n|}{\sqrt{N}} \exp \left( \beta_2 \sqrt{\log N \cdot \log \log N} \right) .$$

but with a better (maybe optimal) parameter  $\beta_2$ .

## 8. Proof idea: upper bound via metric entropy

Step 1: Harald Bohr's reduction: from a random polynomial on the line to a random function on the high-dimensional torus. Let  $\mathbb{S}$  be the unit circle,

$$\pi(N) = |\{p : p \leq N, p \text{ is prime}\}|.$$

Consider the time change application  $f$ ,

$$f : \mathbb{R} \rightarrow \mathbb{S}^{\pi(N)}, \quad f(t) = \left( p_1^{it}, \dots, p_{\pi(N)}^{it} \right).$$

Then the image  $f(\mathbb{R})$  is dense in  $\mathbb{S}^{\pi(N)}$ .

We use the canonical representation,

$$n = \prod_{j \leq \pi(N)} p_j^{b_j(n)}.$$

and obtain for any polynomial

$$\begin{aligned}
 \mathcal{P}_N(t) &= \sum_{n=2}^N a_n n^{it} = \sum_{n=2}^N a_n \left( \prod_{j \leq \pi(N)} p_j^{b_j(n)} \right)^{it} \\
 &= \sum_{n=2}^N a_n \prod_{j \leq \pi(N)} \left( p_j^{b_j(n)} \right)^{it} \\
 &= \tilde{\mathcal{P}}(f(t)),
 \end{aligned}$$

where

$$\tilde{\mathcal{P}}(z) = \sum_{n=2}^N a_n \prod_{j \leq \pi(N)} z_j^{b_j(n)}.$$

Hence,

$$\sup_{t \in \mathbb{R}} |\mathcal{P}(t)| = \sup_{z \in \mathbb{S}^{\pi(N)}} |\tilde{\mathcal{P}}(z)|.$$

Step 2 : choose a clever covering of the torus  $\mathbb{S}^{\pi(N)}$  and use Dudley estimate....

## 9. Lower bound: restriction of the process

First, restriction the of parameter set. We have to evaluate

$$\begin{aligned}
 & \mathbb{E} \sup_{z \in \mathbb{S}^{\pi(N)}} \left| \sum_{n \in \mathcal{V}_{\tau, N}} \varepsilon_n n^{-\sigma} \prod_{j \leq \pi(N)} z_j^{b_j(n)} \right| \\
 = & \mathbb{E} \sup_{z \in \mathbb{S}^{\tau}} \left| \sum_{n \in \mathcal{V}_{\tau, N}} \varepsilon_n n^{-\sigma} \prod_{j \leq \tau} z_j^{b_j(n)} \right|, \quad \varepsilon_n = \pm 1. \\
 & \geq \mathbb{E} \sup_{z \in \mathcal{Z}} \left| \sum_{n \in \mathcal{V}_{\tau, N}} \varepsilon_n n^{-\sigma} \prod_{\tau/2 \leq j \leq \tau} z_j^{b_j(n)} \right|,
 \end{aligned}$$

where

$$\mathcal{Z} = \{z \in \mathbb{S}^{\tau} : z_j = 1, j \leq \tau/2; z_j = \pm 1, j > \tau/2\}.$$

Next, restriction of the process: keep only special terms. Let for  $\tau/2 < j \leq \tau$ ,

$$L_j = \{n \in \mathcal{V}_{\tau, N} : n = p_j \tilde{n}, \quad P_+(\tilde{n}) \leq p_{\tau/2}\}.$$

For  $n \in L_j$  we have  $b_j(n) = 1$  and  $b_i(n) = 0$  elsewhere. We have

$$\begin{aligned} & \mathbb{E} \sup_{z \in \mathcal{Z}} \left| \sum_{n \in \mathcal{V}_{\tau, N}} \varepsilon_n n^{-\sigma} \prod_{\tau/2 \leq j \leq \tau} z_j^{b_j(n)} \right| \\ & \geq \mathbb{E} \sup_{z \in \mathcal{Z}} \left| \sum_{\tau/2 < j \leq \tau} \sum_{n \in L_j} \varepsilon_n n^{-\sigma} \prod_{\tau/2 \leq j \leq \tau} z_j^{b_j(n)} \right| \\ & = \mathbb{E} \sup_{z \in \mathcal{Z}} \left| \sum_{\tau/2 < j \leq \tau} \sum_{n \in L_j} \varepsilon_n n^{-\sigma} z_j \right| \\ & = \sum_{\tau/2 < j \leq \tau} \mathbb{E} \left| \sum_{n \in L_j} \varepsilon_n n^{-\sigma} \right| \approx \sum_{\tau/2 < j \leq \tau} \left| \sum_{n \in L_j} n^{-2\sigma} \right| \\ & \dots \geq \frac{\tau^{1/2} N^{1/2-\sigma}}{(\log \tau)^{1/2}}. \quad \square \end{aligned}$$