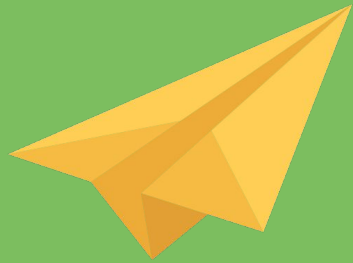
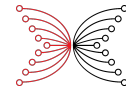


Foundations and Applications of Zero-Knowledge Proofs
4 Sept 2024



zkSNARKs for Blockchains

Anca Nitulescu

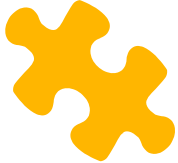


INPUT | OUTPUT

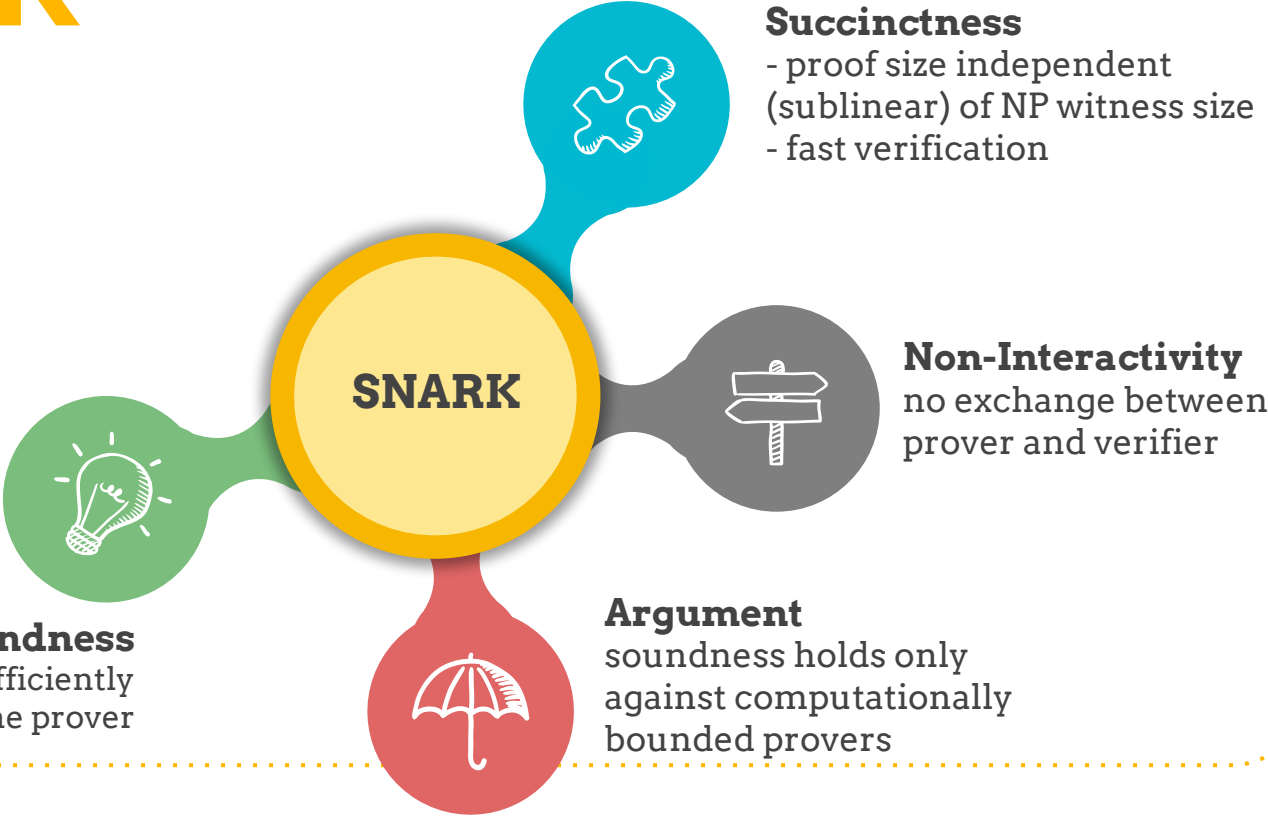


SNARKs

Background



SNARK





zk-SNARK

Zero-Knowledge
does not leak anything
about the witness



Succinctness
- proof size independent
(sublinear) of NP witness size
- fast verification



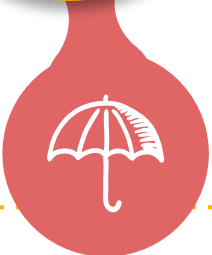
Non-Interactivity
no exchange between
prover and verifier



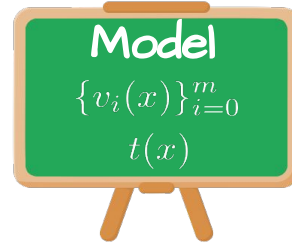
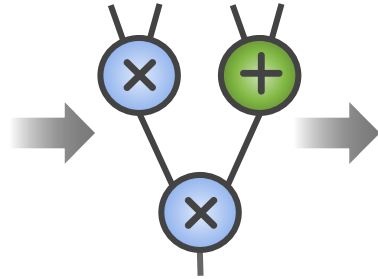
Knowledge Soundness
a witness can be efficiently
extracted from the prover



Argument
soundness holds only
against computationally
bounded provers



Building SNARKs



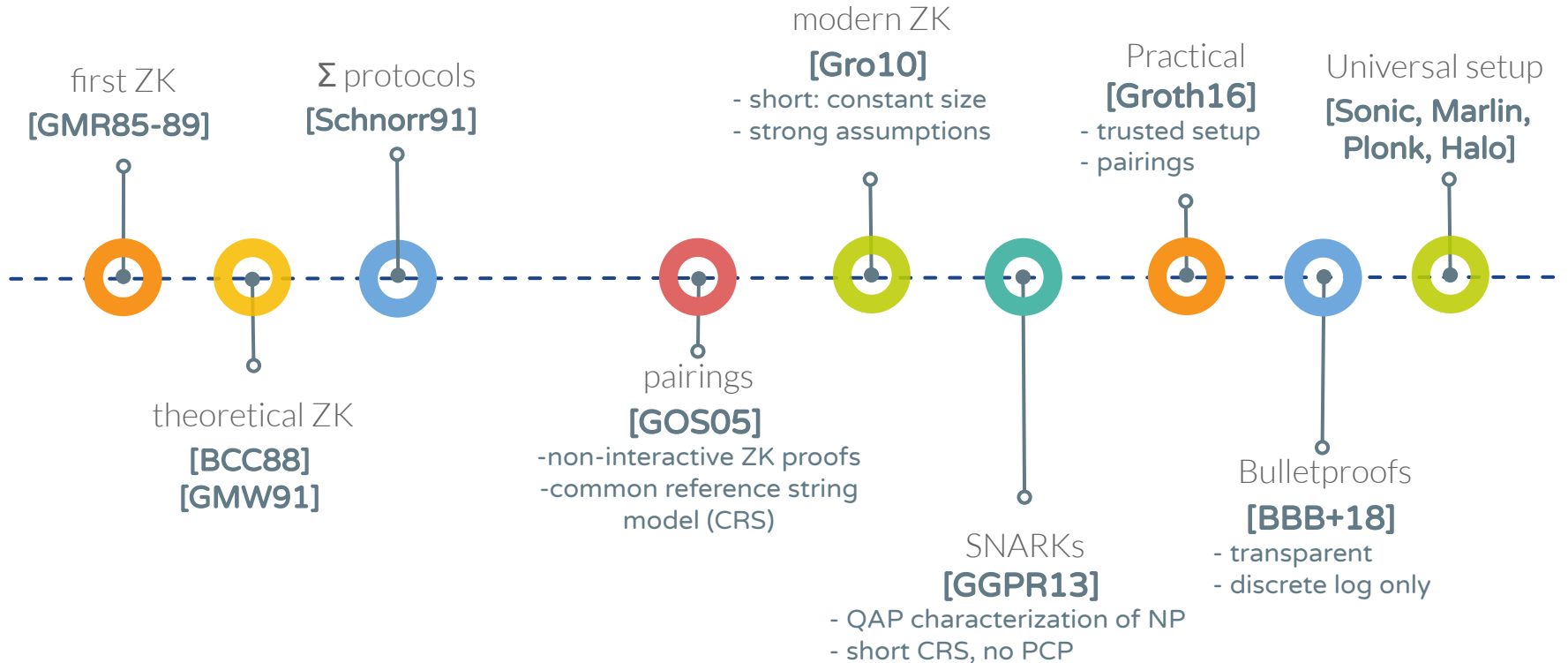
Target Statement
 $\mathcal{L}, R(y,w)=1$

Circuit SAT
NP-complete

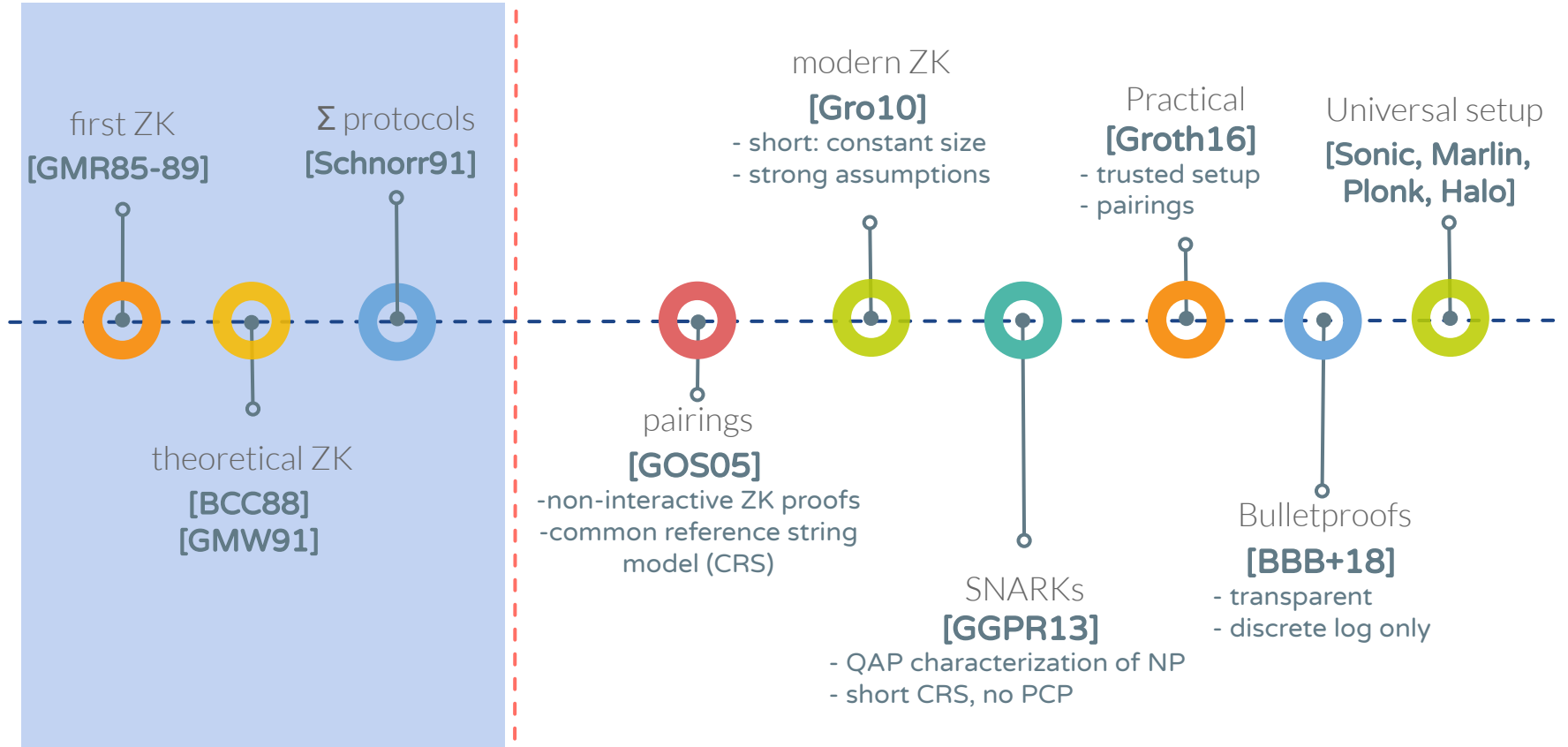
Computational
Model
(Polynomials)

Cryptographic
Compilation

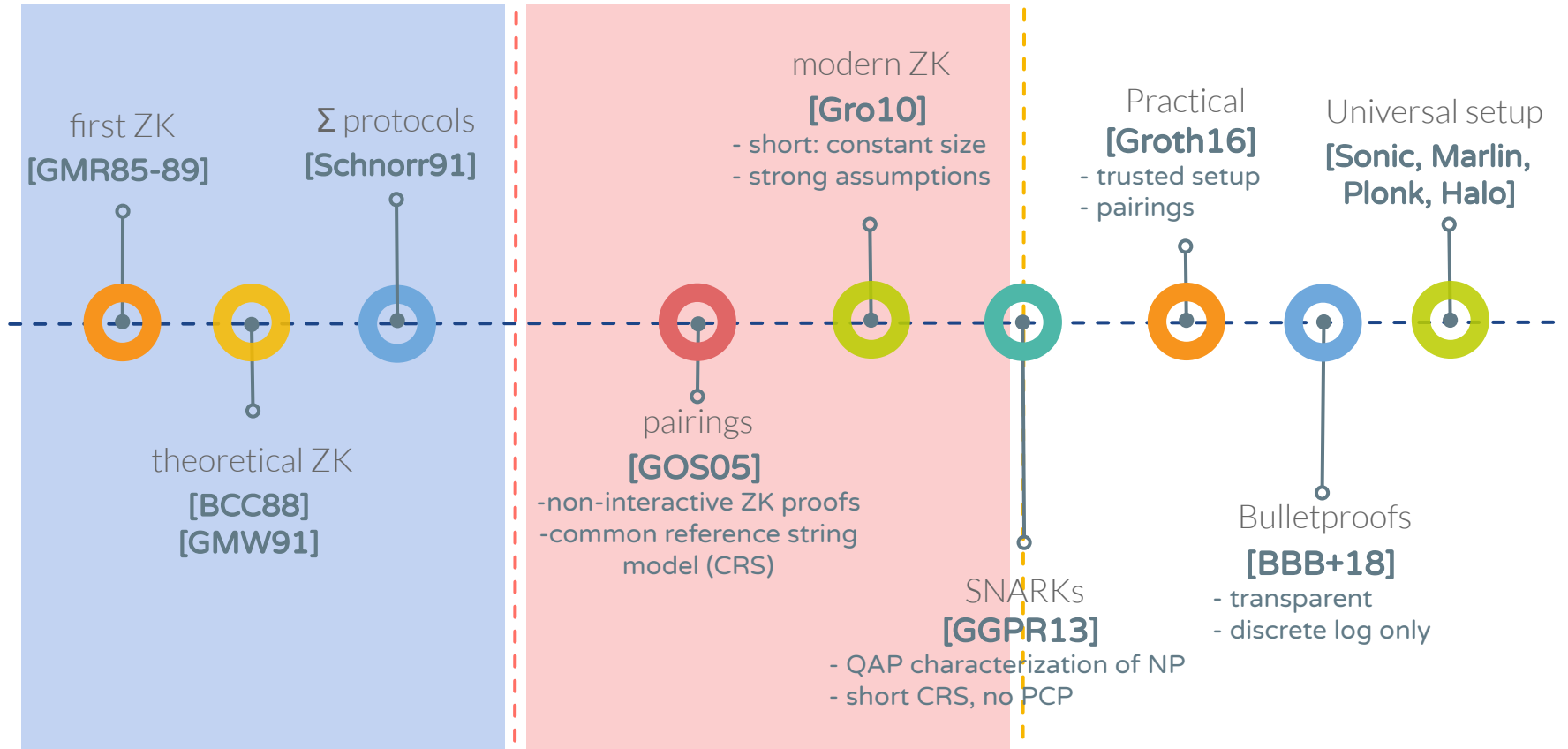
Non-Interactive ZK Proofs



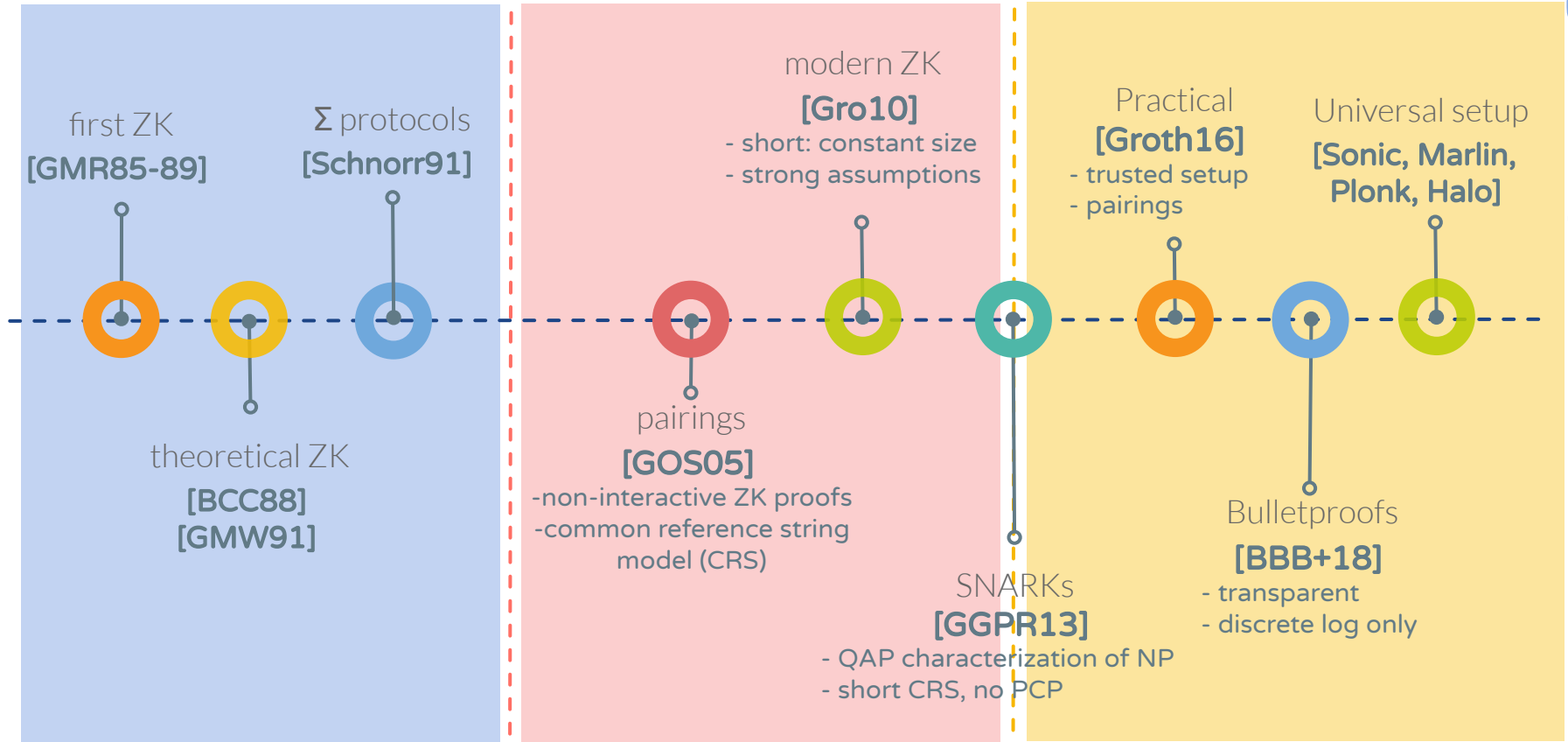
Non-Interactive ZK Proofs



Non-Interactive ZK Proofs



Non-Interactive ZK Proofs



Recent zk-SNARK research

Practical
[Groth16]

- trusted setup
- pairings



Universal setup
**[Sonic, Marlin,
Plonk, Halo]**



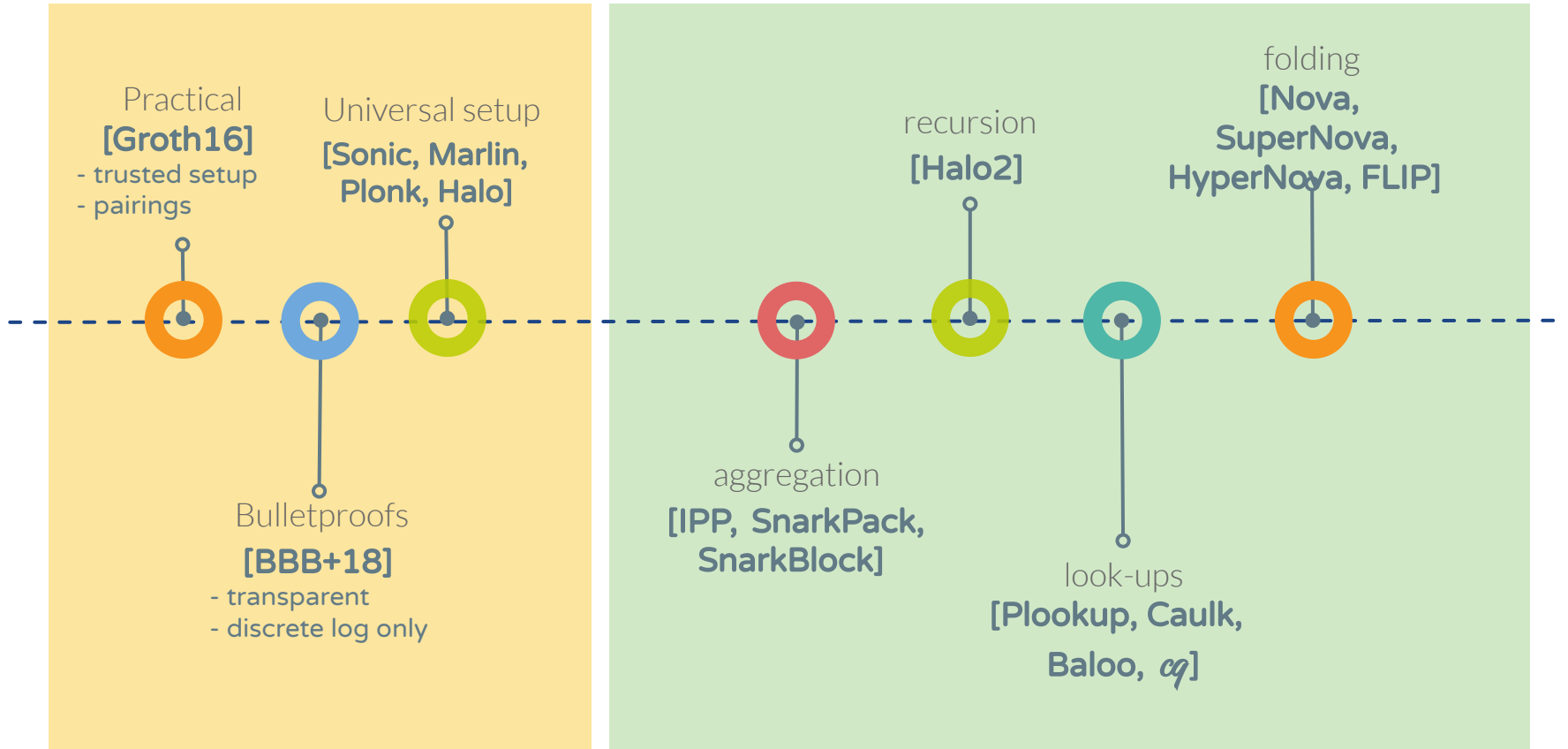
Bulletproofs

[BBB+18]

- transparent
- discrete log only



Recent zk-SNARK research





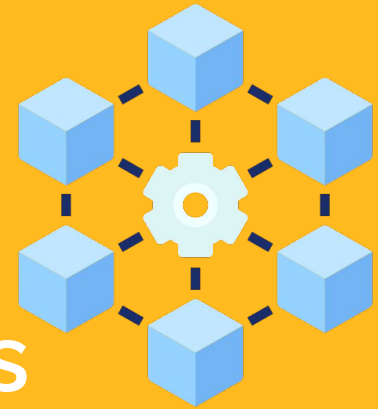
References

- [GMR85-89] **The Knowledge Complexity Of Interactive Proof Systems**, by Shafi Goldwasser, Silvio Micali, and Charles Rackoff
- [GMW91] **How to prove all NP-statements in zero-knowledge, and a methodology of cryptographic protocol design**, by Oded Goldreich, Silvio Micali, and Avi Wigderson
- [BCC88] **Minimum Disclosure Proofs of Knowledge**, by Gilles Brassard, David Chaum, and Claude Crépeau
- [Sch91] **Efficient signature generation by smart cards**, by Claus-Peter Schnorr, Journal of Cryptology, 1991.
- [GOS06] **Perfect Non-interactive Zero Knowledge for NP**, Jens Groth, Rafail Ostrovsky, and Amit Sahai, EUROCRYPT 2006.
- [Groth10] **Short pairing-based non-interactive zero-knowledge arguments**, Jens Groth, ASIACRYPT 2010.
- [KZG10] **Constant-Size Commitments to Polynomials and Their Applications**, by A. Kate, G. Zaverucha, I. Goldberg, in ASIACRYPT'10.
- [GGPRR13] **Quadratic span programs and succinct NIZKs without PCPs**, by R. Gennaro, C. Gentry, Bryan Parno, and Mariana Raykova.
- [BBB+18] **Bulletproofs: Short Proofs for Confidential Transactions and More**, by B. Bünz, J. Bootle, D. Boneh, A. Poelstra, P. Wuille, and G. Maxwell
- [Groth16] **On the Size of Pairing-based Non-interactive Arguments**, Jens Groth, EUROCRYPT 2016.
- [GWC19] **PlonK: Permutations over Lagrange-bases for Oecumenical Noninteractive arguments of Knowledge**, Ariel Gabizon, Zachary J. Williamson, Oana Ciobotaru



Applications

to decentralized systems



Crypto currencies



Distributed Ledger

Alice: 20 Coins
Bob: 55 Coins
Carol: 0 Coins
...
Zoe: 17 Coins



tx: Alice → Bob, 5 Coins



Crypto currencies



Distributed Ledger

Alice: 20 Coins
Bob: 55 Coins
Carol: 0 Coins
...
Zoe: 17 Coins



tx: Alice → Bob, 5 Coins



New State

Alice: 15 Coins
Bob: 60 Coins
Carol: 0 Coins
...
Zoe: 17 Coins

Crypto currencies

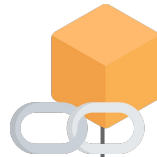


Distributed Ledger

Alice: 20 Coins
Bob: 55 Coins
Carol: 0 Coins
...
Zoe: 17 Coins



tx: Alice → Bob, 5 Coins



- inclusive accountability: anyone can see and check validity of actions
- all transactions are public
- consensus: multiple nodes check a new block



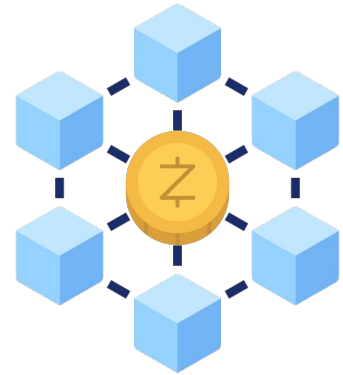
ZK Proofs in Distributed Systems

Privacy:

- Cryptocurrencies that guarantee privacy TX on public Ledger
- **Zero Knowledge Proof:** the TX is valid without revealing more
 - ZCash, Aleo, Aztec, Midnight

Compliance:

- Prove that a private TX is compliant with banking laws
 - Espresso Systems
- Prove that an exchange is solvent
e.g. a company has more assets than obligations to its customers
- **Zero-Knowledge** to protect customers private data





ZK Proofs in Distributed Systems

Distributed storage:

- Peer-to-peer network for storage
- Allows anyone to store and retrieve data on the internet
- built-in economic incentives ensure that files are stored reliably and continuously
- **Proof of Knowledge** of the stored file
 - Filecoin





ZK Proofs in Distributed Systems

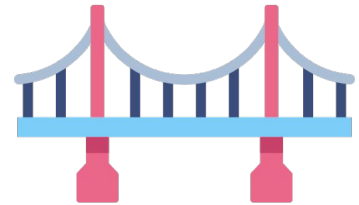
Distributed storage:

- Peer-to-peer network for storage
- Allows anyone to store and retrieve data on the internet
- built-in economic incentives ensure that files are stored reliably and continuously
- **Proof of Knowledge** of the stored file
 - Filecoin



Bridging Blockchains:

- Transfer of assets from a chain to another
- Source chain proves to target chain that the asset is locked-up
- No need for Zero Knowledge
- Need for **Proof of Knowledge**





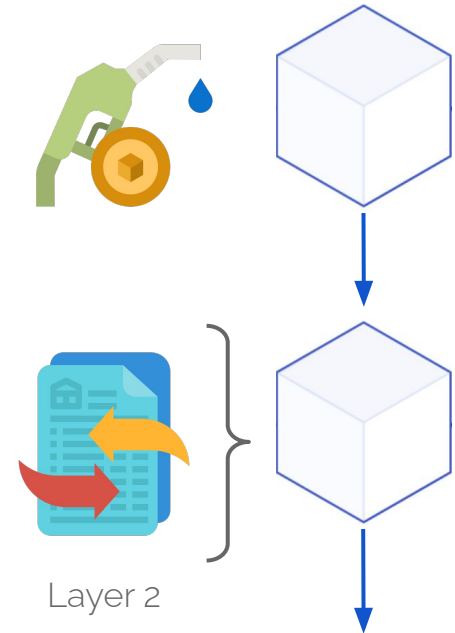
ZK Proofs in Distributed Systems

Rollups = Outsourced Computation

- Scalability tool
- Off-chain services (Layer 2) process a batch of TX (validates TX) and proves the work
- Main chain (Layer 1) verifies efficiently the work of an off-chain service
- Reduces the cost of transactions
- **No need** for Zero Knowledge
- Need for **Succinctness & Proof of Knowledge**

No zk: ZKSync, Loopring, Linea by Consensys, ZKSpace

With zk: Aztec, Polygon Hermez





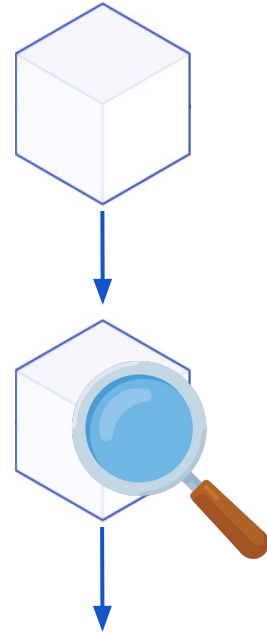
ZK Proofs Key Requirements

Key Properties for Proofs in Blockchains

- non-interactivity
- publicly verifiable
- **succinctness for proofs**
- **efficiency for verifier**

Answer: **SNARKs**

Succinct Non-Interactive Arguments of Knowledge





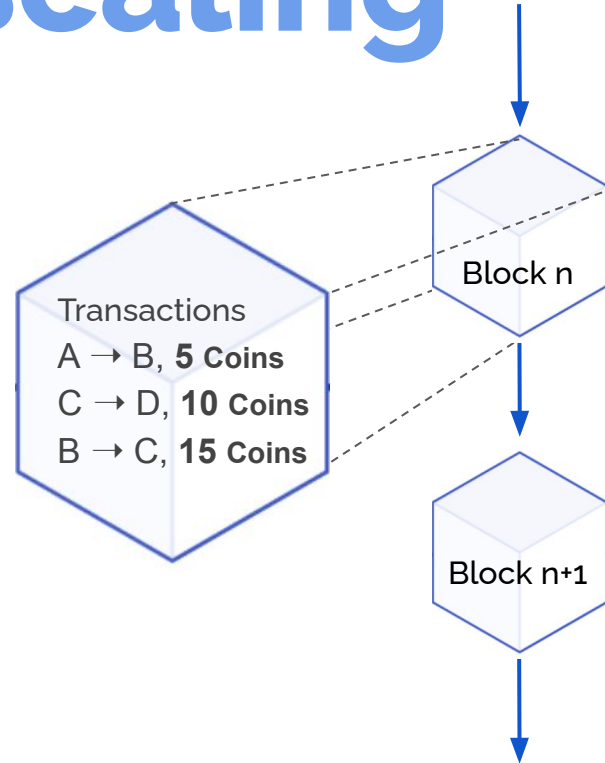
What is left?

Blockchain scaling



ZK Roll-ups

- instead of posting all transaction data on-chain
- execute **tx** using off-chain computation
- a summary of the state changes is published to the chain and verified



Blockchain scaling

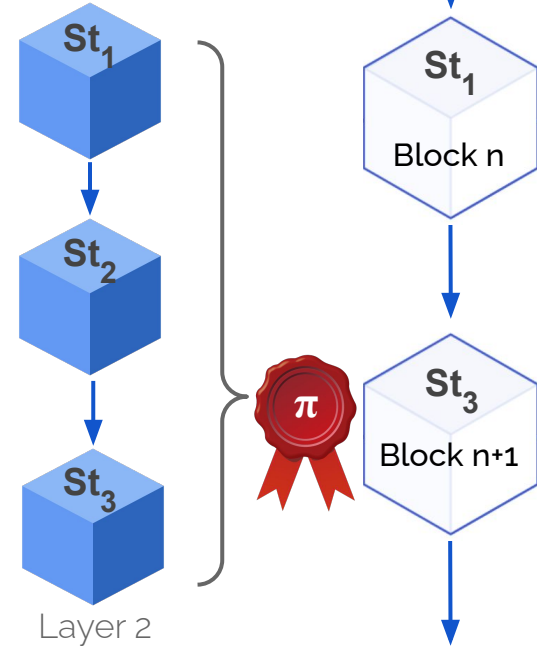


ZK Roll-ups

- instead of posting all transaction data on-chain
- execute **tx** using off-chain computation
- a summary of the state changes is published to the chain and verified

SNARK π proves the correctness of the changes

Statement: state changes proposed by the layer 2 are correct,
i.e. are the result of the execution of the given batch of transactions



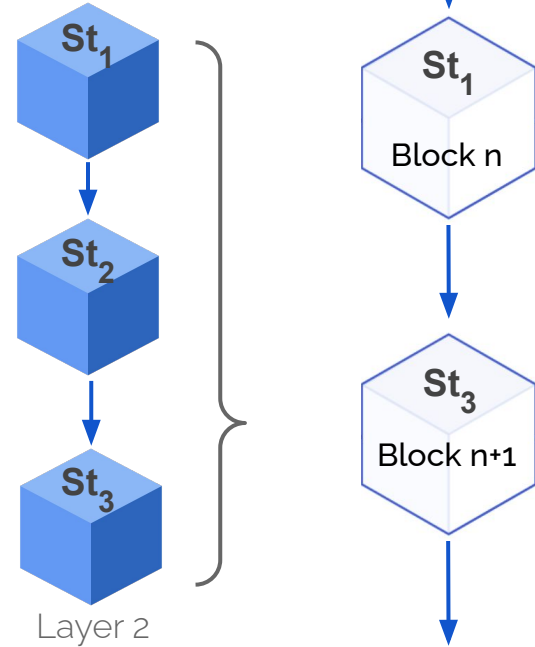
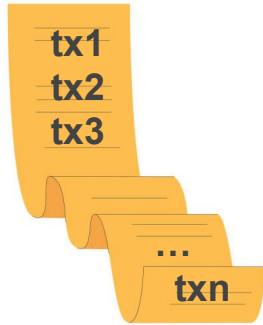
Blockchain scaling



ZK Roll-ups

Prove multiple instances:

- correctness of the execution of batch of **tx**



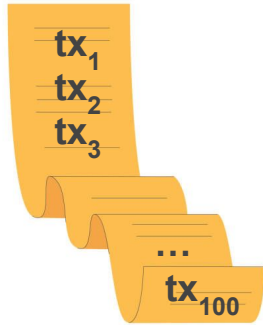
Blockchain scaling



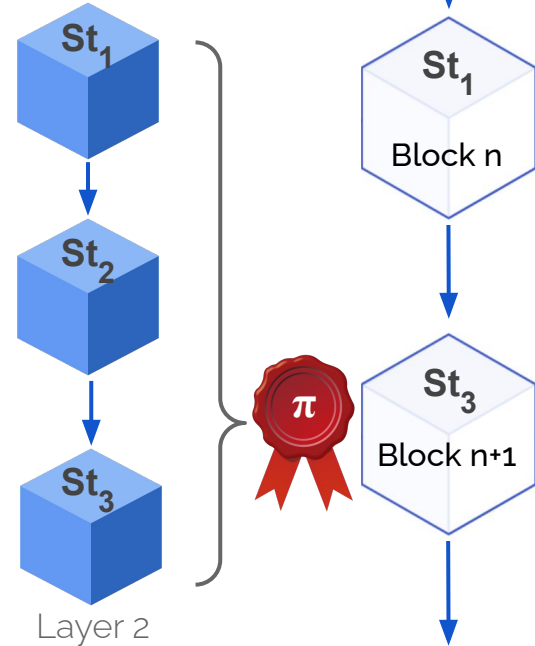
ZK Roll-ups

Example 100 Tx

Naively: generate state transition proof once all 100 Tx are submitted → **long delay**



Proving is slow

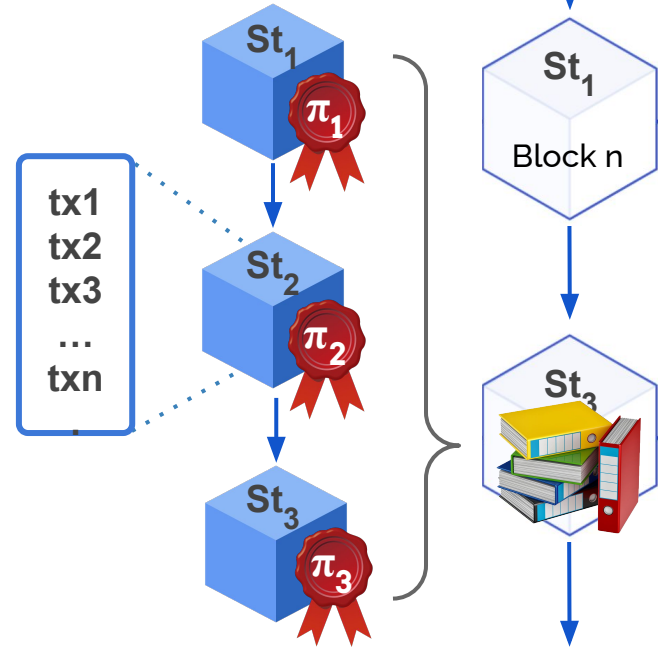


Blockchain scaling



ZK Roll-ups

Periodically provide proofs for valid transaction batches

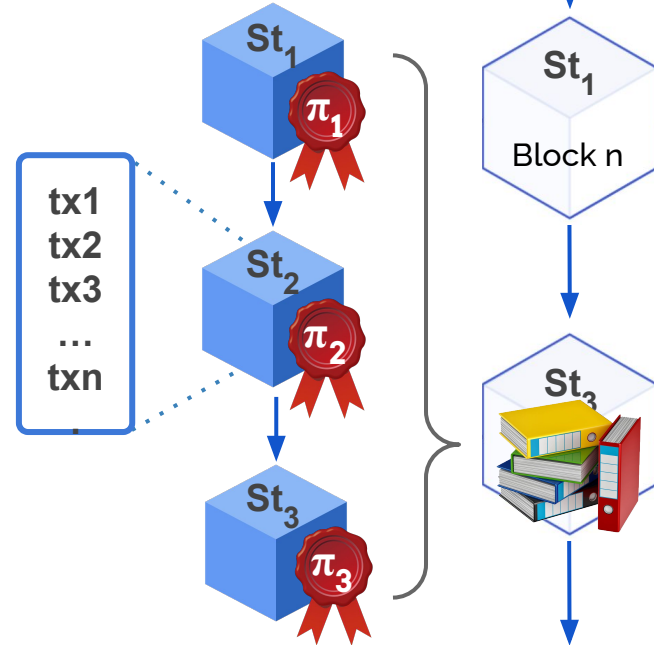
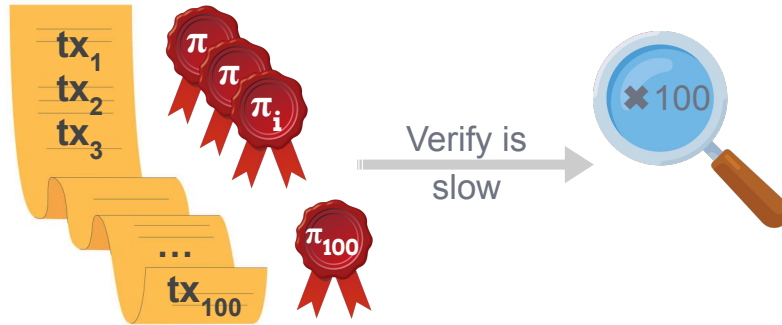


Blockchain scaling

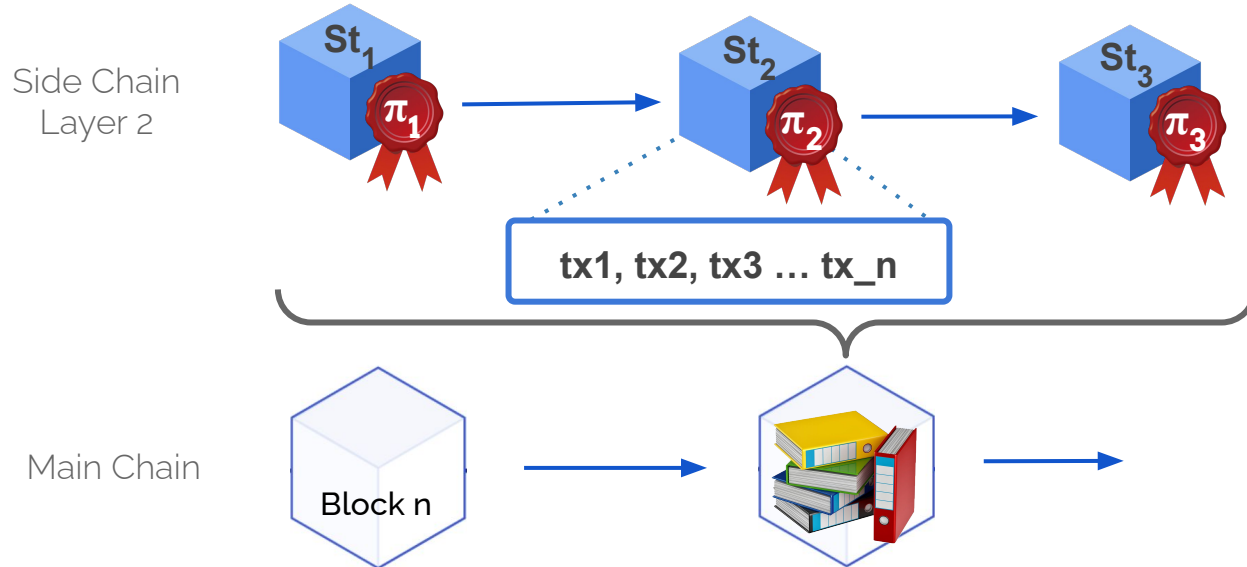


ZK Roll-ups

Periodically provide proofs for valid transaction batches

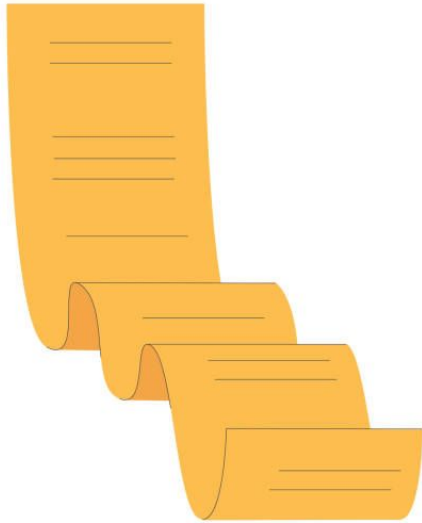


Prove & scale ?





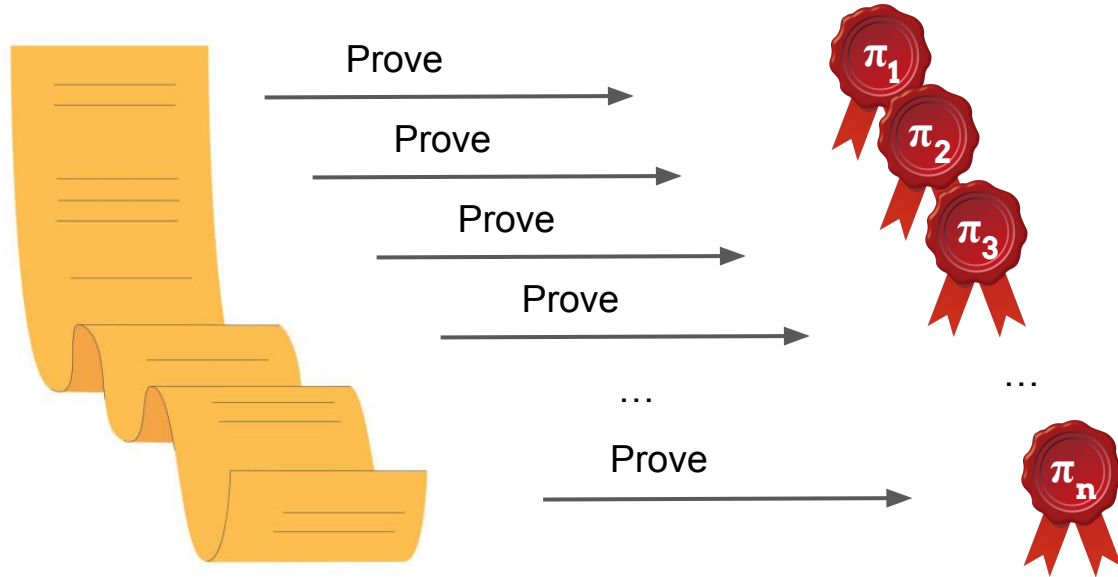
In Brief



Multiple instances to prove



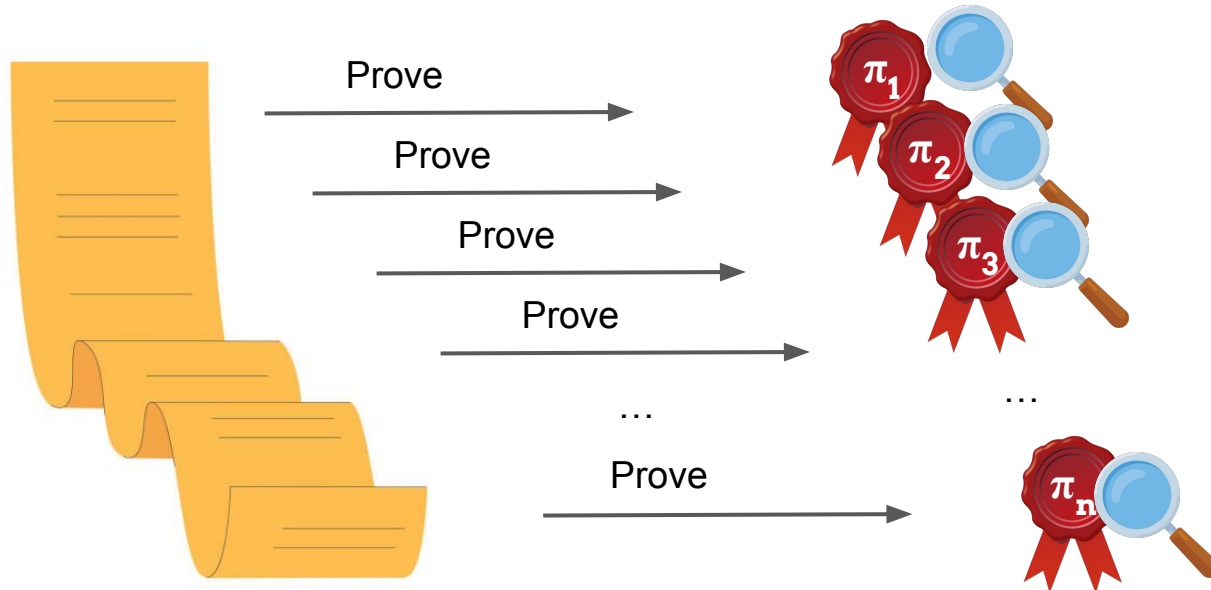
Naively



Multiple instances to prove



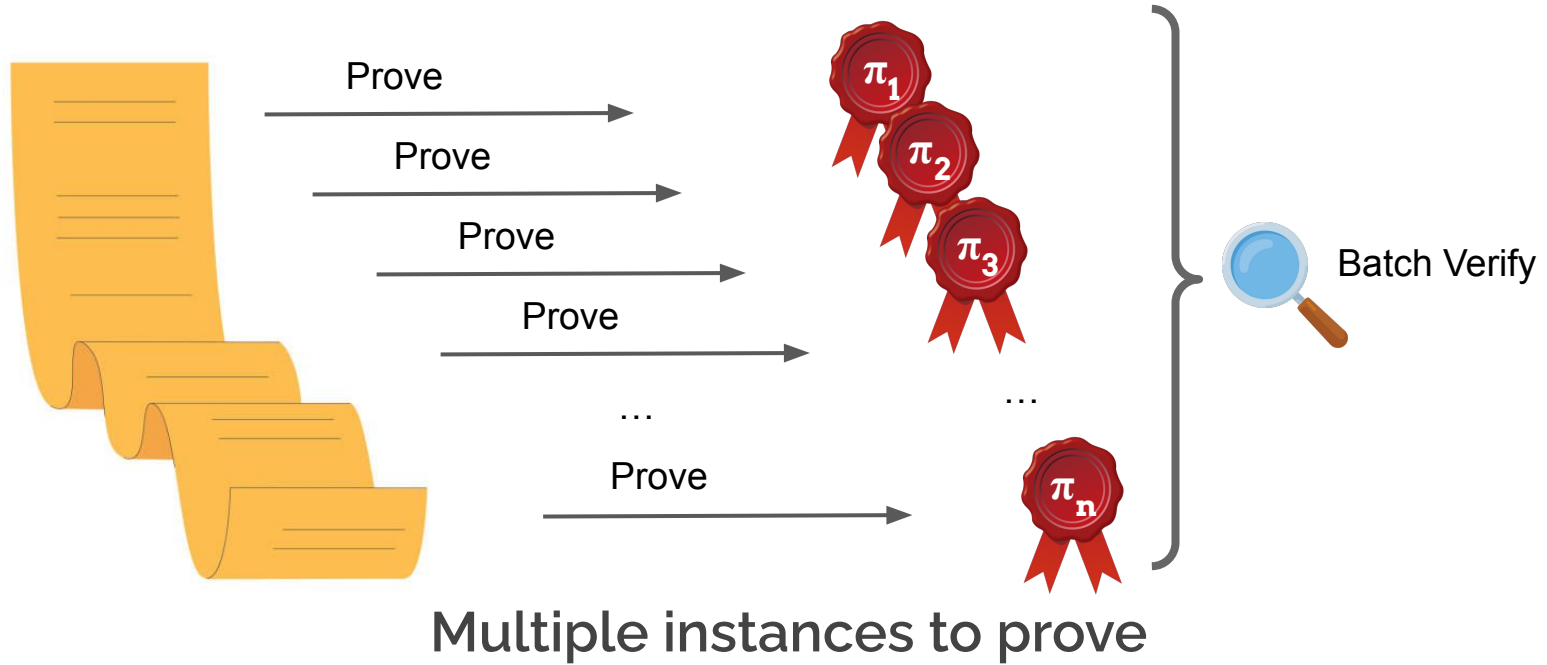
Naively



Multiple instances to prove

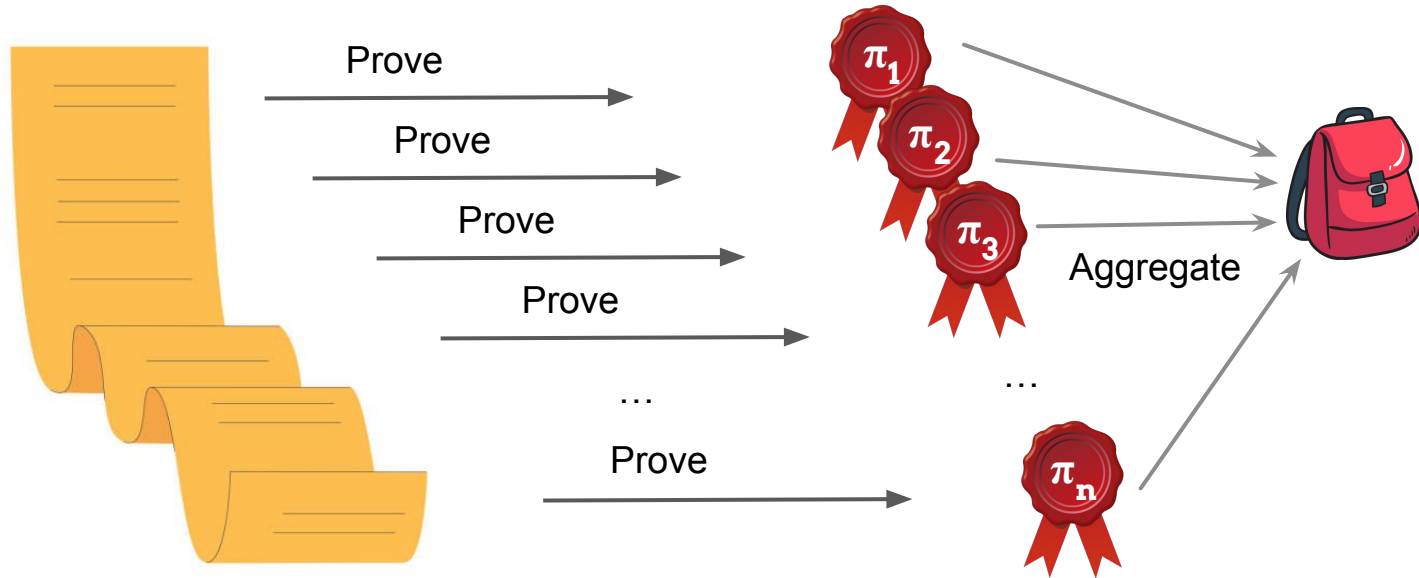


Batch Verify





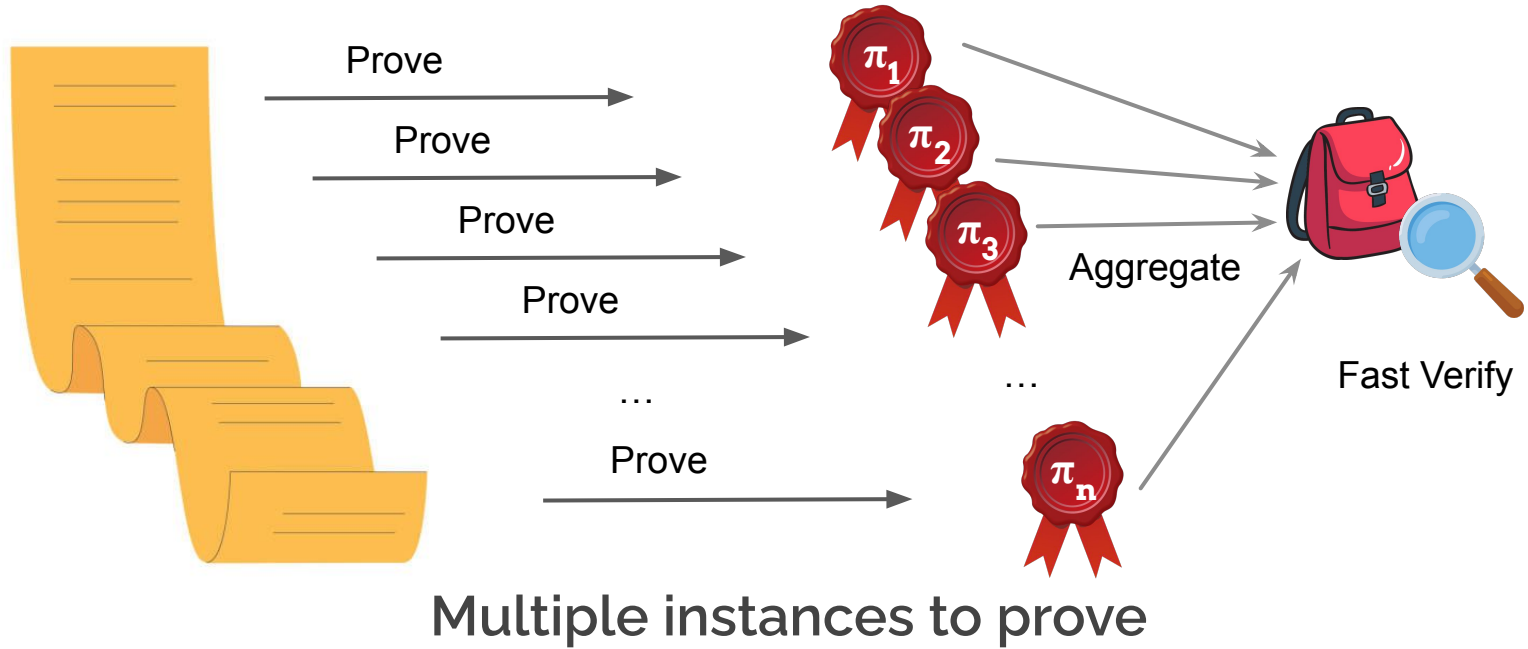
Aggregation



Multiple instances to prove

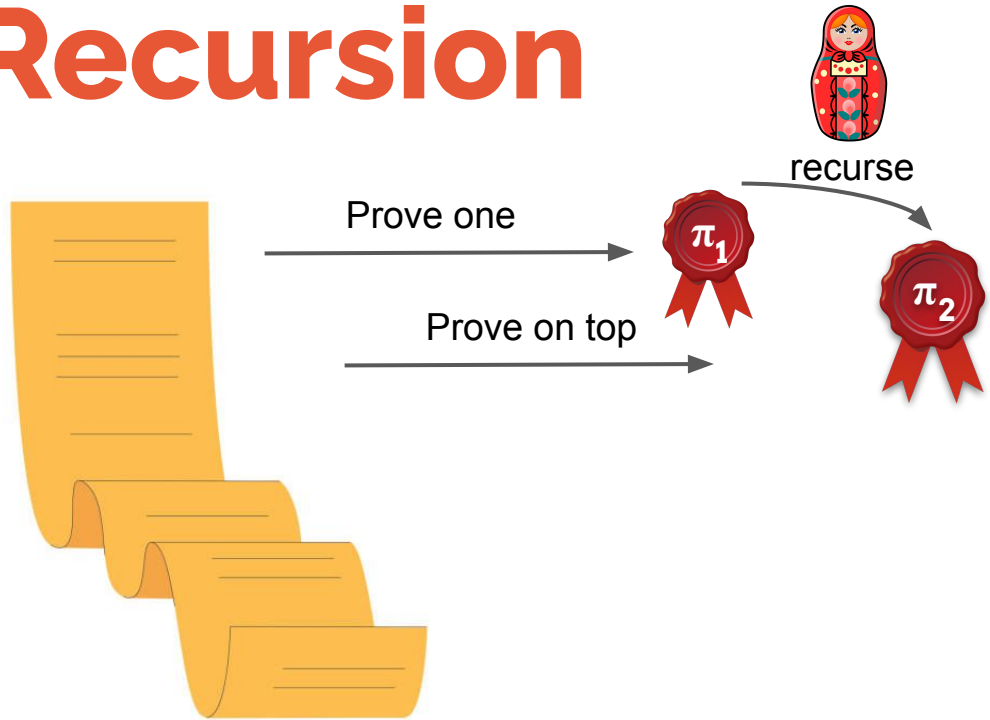


Aggregation





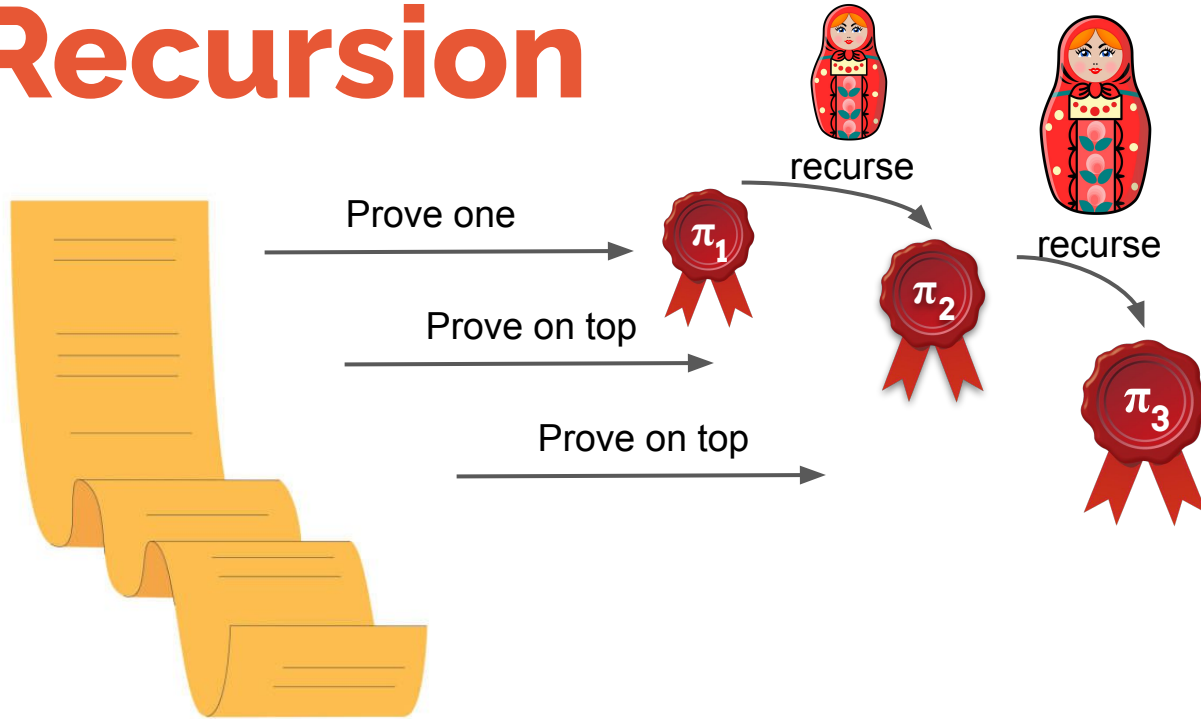
Recursion



Multiple instances to prove



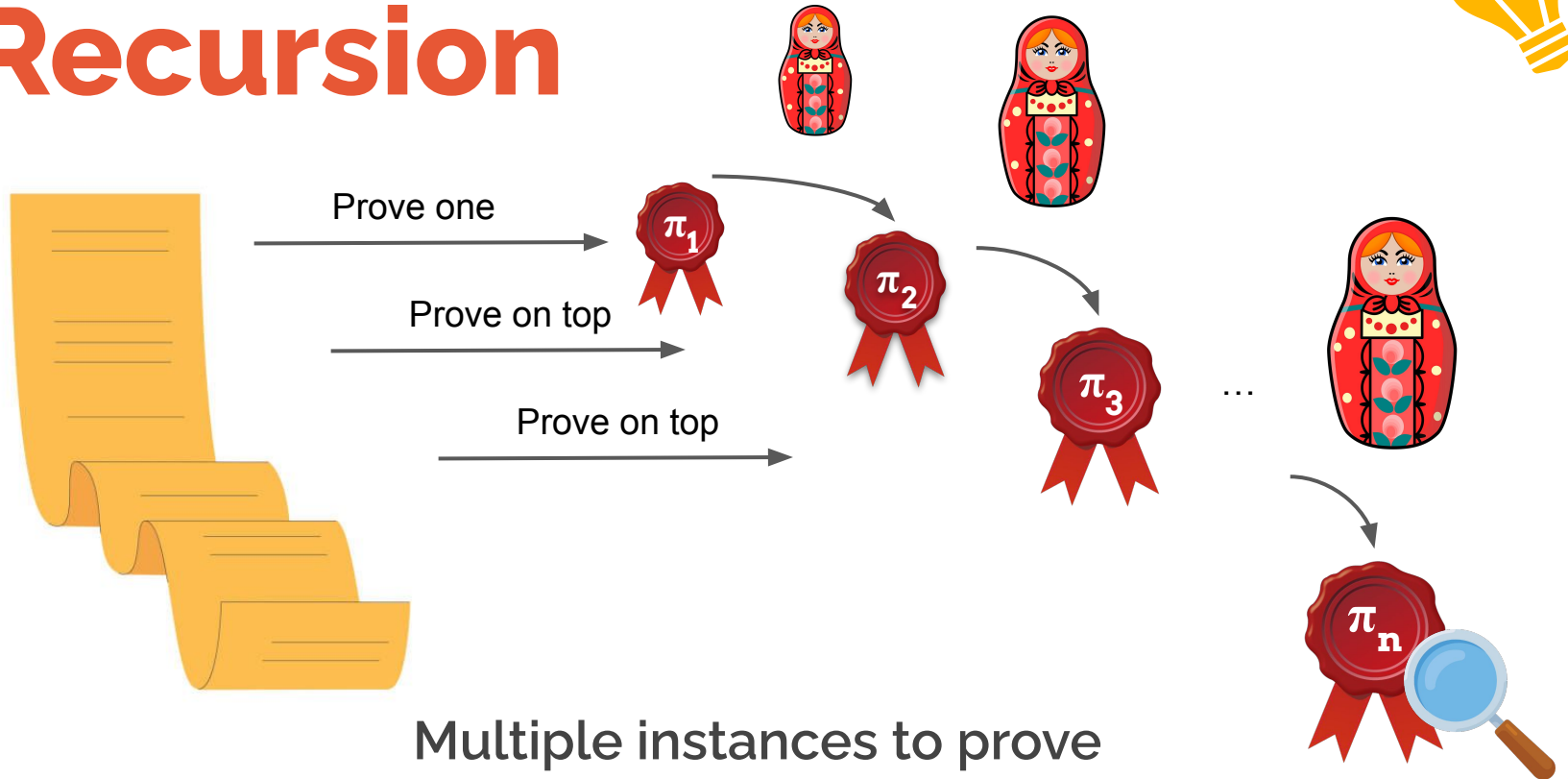
Recursion



Multiple instances to prove



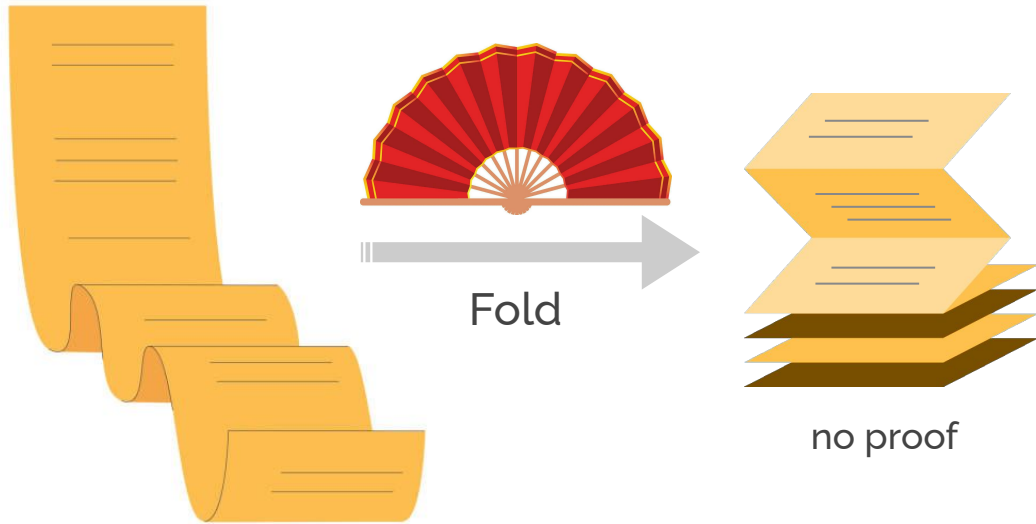
Recursion



Multiple instances to prove



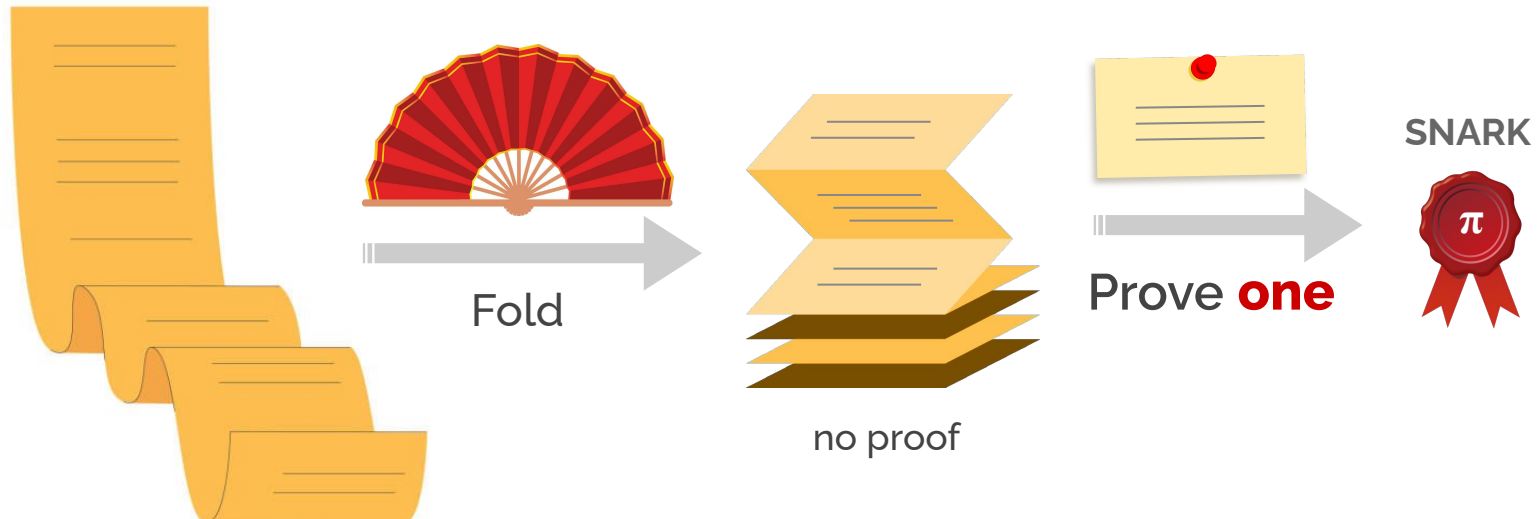
Folding



Multiple instances to prove



Folding



Multiple instances to prove



Prove many instances



Aggregation

Proof Size



Verifier Time



Prover Time



Recursion



Folding





SNARK Aggregation



Prove many instances



Aggregation

Proof Size



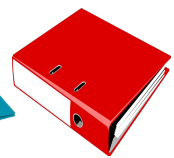
Verifier Time



Prover Time



Recursion

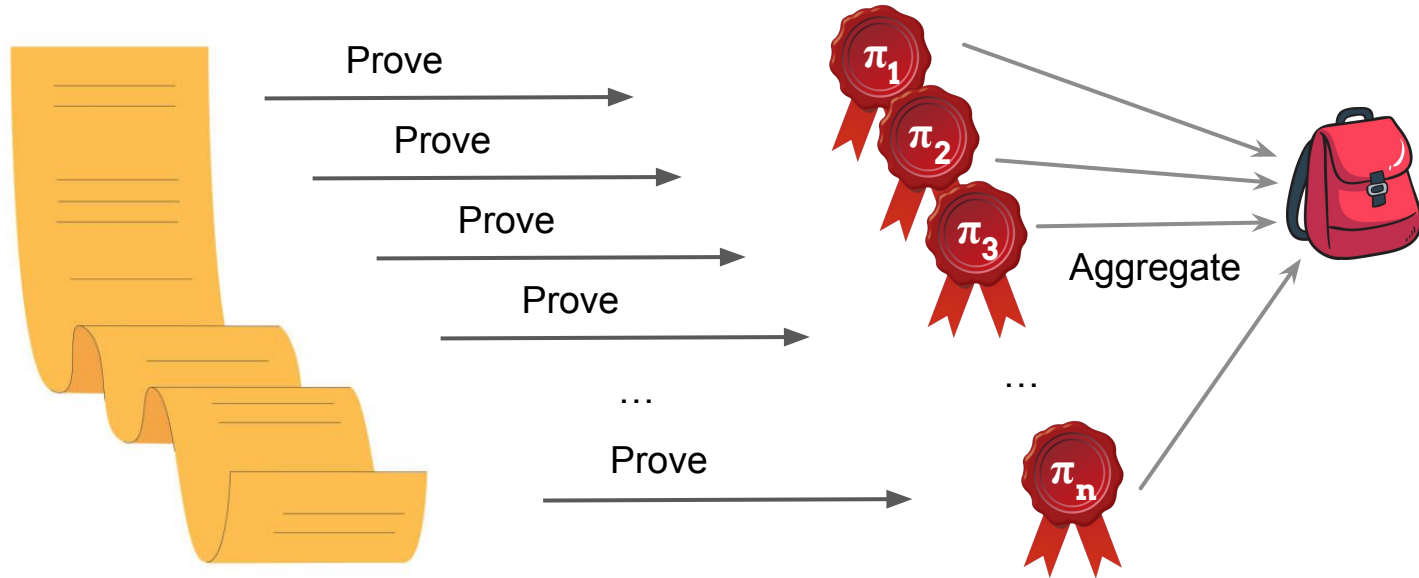


Folding





Aggregation



Multiple instances to prove

Distributed storage



Storage Providers

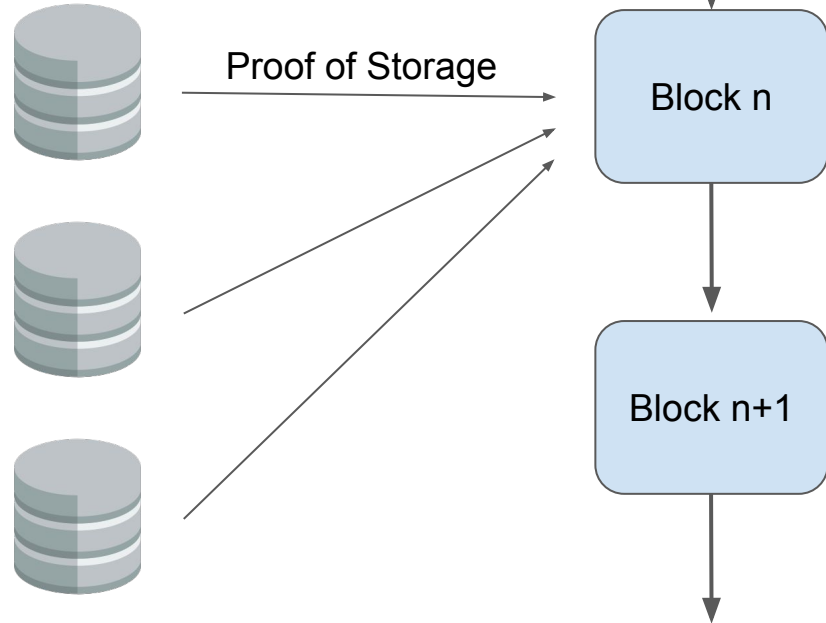
- onboard storage capacity
- earn block rewards
- regularly prove the storage

= **Provers**

Nodes in network

- ensure data is being stored, maintained, and secured
- need to check proofs of space

= **Verifiers**





Proof of storage



unit1

Proves 32GB



submit on chain



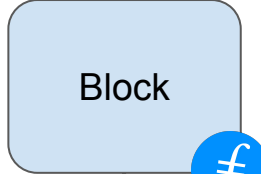
unit2

Proves 32GB



unit3

Proves 32GB



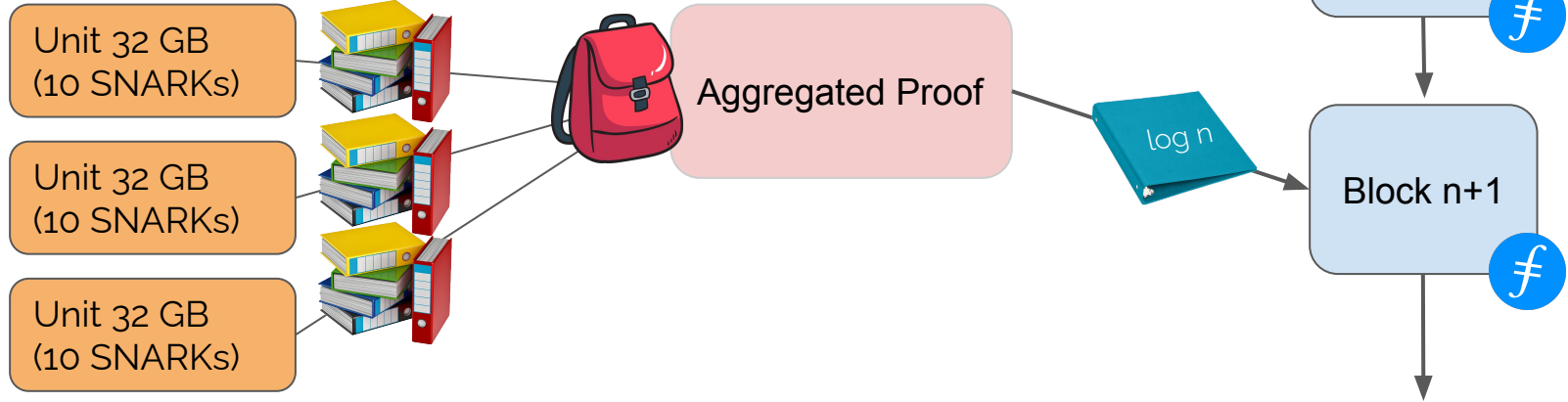
=> storage onboarding limit



Snark Pack Aggregation

SnarkPack: Practical SNARK Aggregation Nicolas Gailly, Mary Maller, Anca Nitulescu

- Size is **logarithmic** as well as verification time
- Prover overhead





Bilinear Groups

Bilinear Groups

$$[a]_1 \in \mathbf{G}_1, [b]_2 \in \mathbf{G}_2$$

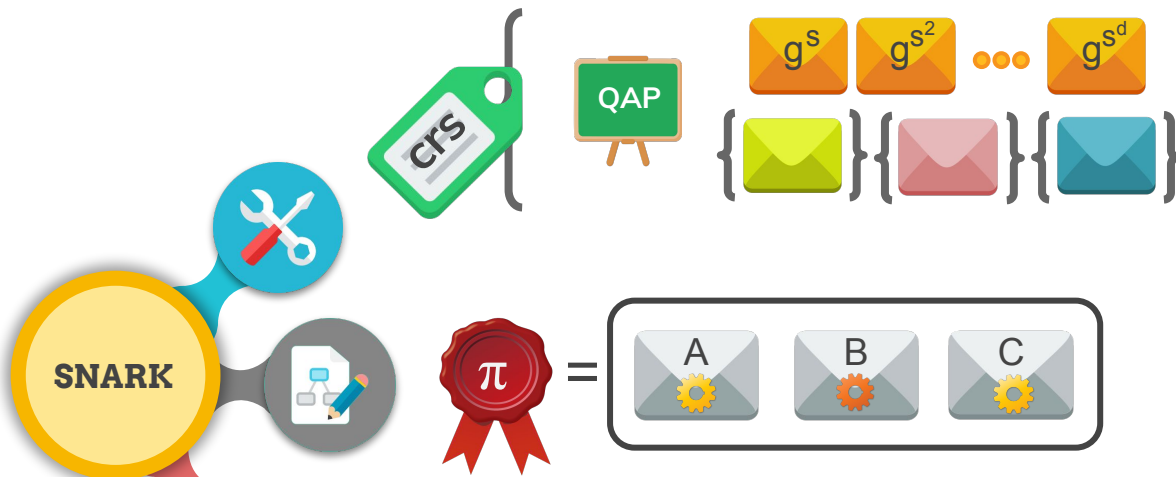
$$e : \mathbf{G}_1 \times \mathbf{G}_2 \rightarrow \mathbf{G}_T$$

$$e([a]_1, [b]_2) = [ab]_T$$





Groth 16



$$e(A, B) = e(C, D)$$

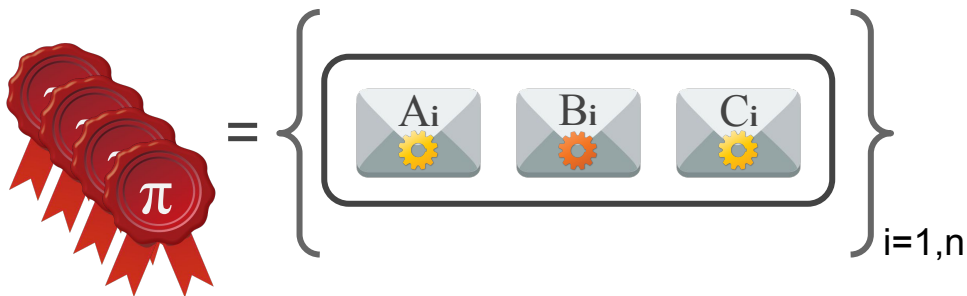


Bilinear Groups
 $[a]_1 \in \mathbf{G}_1, [b]_2 \in \mathbf{G}_2$
 $e : \mathbf{G}_1 \times \mathbf{G}_2 \rightarrow \mathbf{G}_T$
 $e([a]_1, [b]_2) = [ab]_T$

Many SNARKs



Proofs



Verification ($D = g^d$)

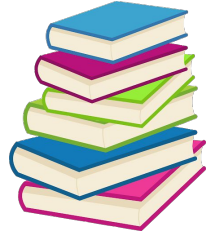
$$e(A_1, B_1) = e(C_1, D)$$

$$e(A_2, B_2) = e(C_2, D)$$

...

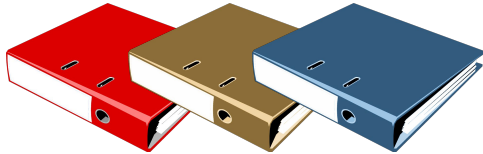
$$e(A_n, B_n) = e(C_n, D)$$

Verify many **SNARKs**



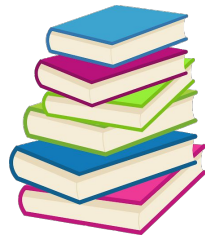
Batch Verification

Proof Size



Verification
Time





SNARK Batching

Verification

$$e(A_1, B_1) = e(C_1, D)$$

$$e(A_2, B_2) = e(C_2, D)$$

...

$$e(A_n, B_n) = e(C_n, D)$$



SNARK Batching

Verification

$$e(A_1, B_1) = e(C_1, D)$$

$$e(A_2, B_2) = e(C_2, D)$$

...

$$e(A_n, B_n) = e(C_n, D)$$

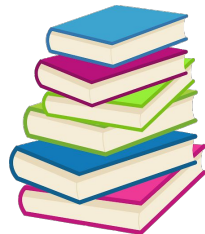


$$r \times e(A_1, B_1) = e(C_1, D)$$

$$r^2 \times e(A_2, B_2) = e(C_2, D)$$

...

$$r^n \times e(A_n, B_n) = e(C_n, D)$$



SNARK Batching

Verification

$$e(A_1, B_1) = e(C_1, D)$$

$$e(A_2, B_2) = e(C_2, D)$$

...

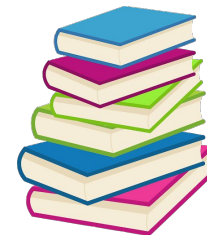
$$e(A_n, B_n) = e(C_n, D)$$



Batch Verification

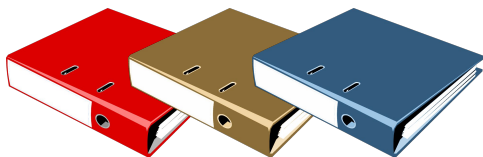
$$\sum r^i e(A_i, B_i) = \sum r^i e(C_i, D)$$

Verify many SNARKs



Batch Verification

Proof Size



Verification Time



Aggregation

Proof Size



Verification Time





SNARK Aggregation

SnarkPack: Practical SNARK Aggregation Nicolas Gailly, Mary Maller, Anca Nitulescu

Batch Verification

$$\sum e(A_i, r^i B_i) = \sum e(r^i C_i, D)$$



$$\sum e(A_i, r^i B_i) = e(\sum r^i C_i, D)$$

Bilinear Groups

$$[a]_1 \in \mathbb{G}_1, [b]_2 \in \mathbb{G}_2$$

$$e : \mathbb{G}_1 \times \mathbb{G}_2 \rightarrow \mathbb{G}_T$$

$$e([a]_1, [b]_1) = [ab]_T$$



SNARK Aggregation

SnarkPack: Practical SNARK Aggregation Nicolas Gailly, Mary Maller, Anca Nitulescu

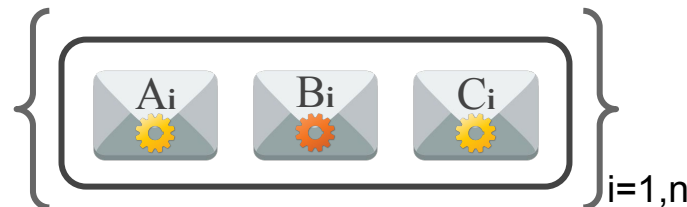
Batch Verification

$$\sum e(A_i, r^i B_i) = \sum e(r^i C_i, D)$$



$$\sum e(A_i, r^i B_i) = e(\sum r^i C_i, D)$$

Aggregation





SNARK Aggregation

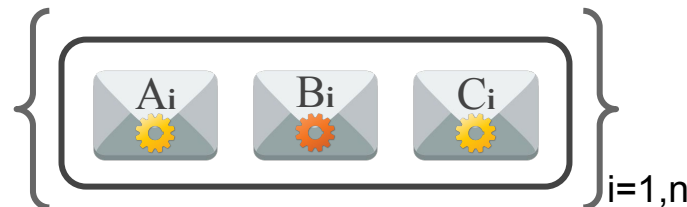
Batch Verification

$$\sum e(A_i, r^i B_i) = \sum e(r^i C_i, D)$$



$$\sum e(A_i, r^i B_i) = e(\sum r^i C_i, D)$$

Aggregation



$$Z_{AB} = \sum e(A_i, r^i B_i)$$

$$Z_C = \sum r^i C_i$$





SNARK Aggregation

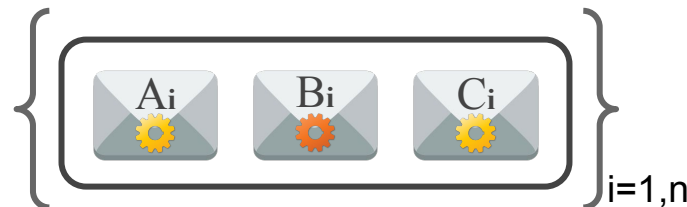
Batch Verification

$$\sum e(A_i, r^i B_i) = \sum e(r^i C_i, D)$$



$$Z_{AB} = e(Z_C, D)$$

Aggregation



$$Z_{AB} = \sum e(A_i, r^i B_i)$$

$$Z_C = \sum r^i C_i$$





SNARK Aggregation

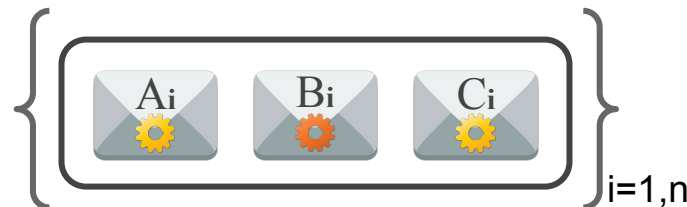
Batch Verification

$$\sum e(A_i, r^i B_i) = \sum e(r^i C_i, D)$$



$$Z_{AB} = e(Z_C, D)$$

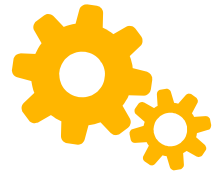
Aggregation



$$Z_{AB} = \sum e(A_i, r^i B_i)$$

$$Z_C = \sum r^i C_i$$





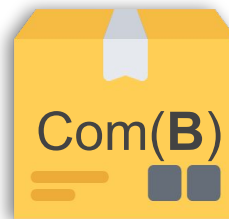
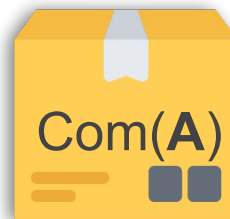
Tools: MIPP & TIPP

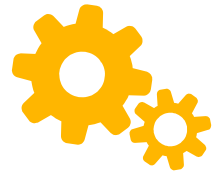
Proofs for Inner Pairing Products and Applications - Bünz, Maller, Mishra, Tyagi, Vesely

$$\langle \mathbf{A}, \mathbf{r} \rangle = \sum r_i A_i$$

$$A_i \in \mathbb{G}_1, B_i \in \mathbb{G}_2, r_i \in \mathbb{Z}_q$$

$$\langle \mathbf{A}, \mathbf{B} \rangle = \sum e(A_i, B_i)$$





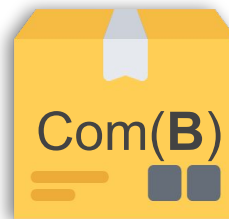
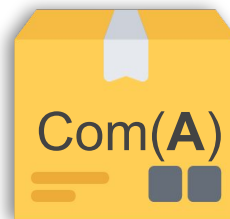
Tools: MIPP & TIPP

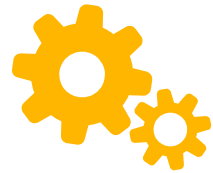
Proofs for Inner Pairing Products and Applications - Bünz, Maller, Mishra, Tyagi, Vesely

$$\langle \mathbf{A}, \mathbf{r} \rangle = \sum r_i A_i = \sum [a_i, r_i]_1$$

$$A_i \in \mathbb{G}_1, B_i \in \mathbb{G}_2, r_i \in \mathbb{Z}_q$$

$$\langle \mathbf{A}, \mathbf{B} \rangle = \sum e(A_i, B_i) = [\sum a_i b_i]_{\mathcal{T}}$$





Tools: MIPP & TIPP

Proofs for Inner Pairing Products and Applications - Bünz, Maller, Mishra, Tyagi, Vesely

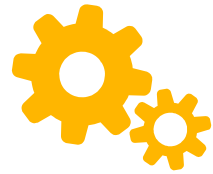
$$\langle \mathbf{C}, \mathbf{r} \rangle = \sum r_i A_i$$

$$\langle \mathbf{A}, \mathbf{rB} \rangle = \sum e(A_i, B_i)$$

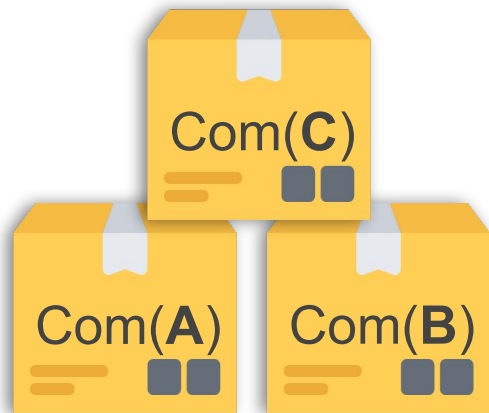
$$Z_{\mathbf{C}} = \langle \mathbf{C}, \mathbf{r} \rangle$$

$$Z_{\mathbf{AB}} = \langle \mathbf{A}, \mathbf{rB} \rangle$$

Aggregation

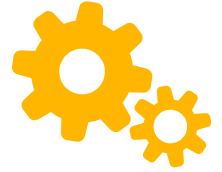


MIPP & TIPP Strategy



$$Z_C = \langle C, r \rangle$$

$$Z_{AB} = \langle A, r_B \rangle$$

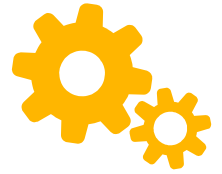


TIPP Folding Strategy

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots, A_n)$$

$$\mathbf{B} = (B_1, B_2, \dots, B_n)$$

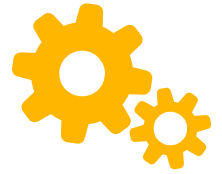


Split & Collapse

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots, A_n)$$

$$\mathbf{B} = (B_1, B_2, \dots, B_n)$$

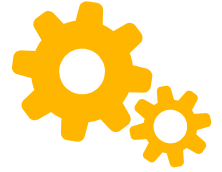


Split & Collapse

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots A_n)$$

$$\mathbf{B} = (B_1, B_2 \dots B_n)$$



Split & Collapse

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots A_n)$$

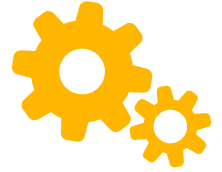
$$\mathbf{B} = (B_1, B_2 \dots B_n)$$

A_{left}

A_{right}

B_{left}

B_{right}



Split & Collapse

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots, A_n)$$

\mathbf{A}_{left}

$\mathbf{A}_{\text{right}}$

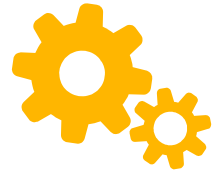
$$\mathbf{A}' = (A'_1, \dots, A'_{n/2})$$

$$\mathbf{B} = (B_1, B_2, \dots, B_n)$$

\mathbf{B}_{left}

$\mathbf{B}_{\text{right}}$

$$\mathbf{B}' = (B'_1, \dots, B'_{n/2})$$



Split & Collapse

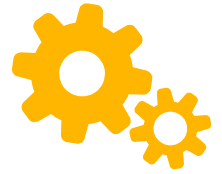
$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

$$\mathbf{A} = (A_1, A_2, \dots, A_n)$$

$$\mathbf{B} = (B_1, B_2, \dots, B_n)$$

$$\mathbf{A}' = (A'_1, \dots, A'_{n/2})$$

$$\mathbf{B}' = (B'_1, \dots, B'_{n/2})$$



Split & Collapse

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(A_1, B_1) + e(A_2, B_2) \dots + e(A_n, B_n)$$

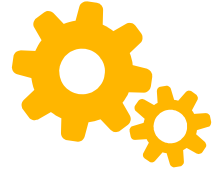
$$\mathbf{A} = (A_1, A_2, \dots, A_n)$$

$$\mathbf{B} = (B_1, B_2, \dots, B_n)$$

$$\langle \mathbf{A}', \mathbf{B}' \rangle$$

$$\mathbf{A}' = (A'_1, \dots, A'_{n/2})$$

$$\mathbf{B}' = (B'_1, \dots, B'_{n/2})$$



Split & Collapse

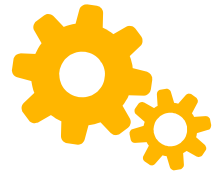
Split

$$\mathbf{A} = (\mathbf{A}_{\text{left}}, \mathbf{A}_{\text{right}})$$
$$\mathbf{B} = (\mathbf{B}_{\text{left}}, \mathbf{B}_{\text{right}})$$
$$\mathbf{Z}'_{\mathbf{A},\mathbf{B}} = \langle \mathbf{A}_{\text{half}}, \mathbf{B}_{\text{half}} \rangle$$

$$\mathbf{L} = \langle \mathbf{A}_{\text{right}}, \mathbf{B}_{\text{left}} \rangle \quad \mathbf{R} = \langle \mathbf{A}_{\text{left}}, \mathbf{B}_{\text{right}} \rangle$$

$$\mathbf{Z}'_{\mathbf{A},\mathbf{B}} = \mathbf{L} + \langle \mathbf{A}, \mathbf{B} \rangle + \mathbf{R}\mathbf{x}$$

$$\langle \mathbf{A}, \mathbf{B} \rangle = e(\mathbf{A}_1, \mathbf{B}_1) + e(\mathbf{A}_2, \mathbf{B}_2) \dots + e(\mathbf{A}_n, \mathbf{B}_n)$$



Linear Verification



$$Z_{AB} = \langle A, B \rangle$$

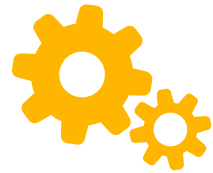
$$A = (A_1, A_2, \dots, A_n)$$



$$A' = A_{\text{left}} A_{\text{right}}$$



$$A_{\text{final}}$$



Linear Verification



$$Z_{AB} = \langle A, B \rangle$$

$$A = (A_1, A_2, \dots, A_n)$$

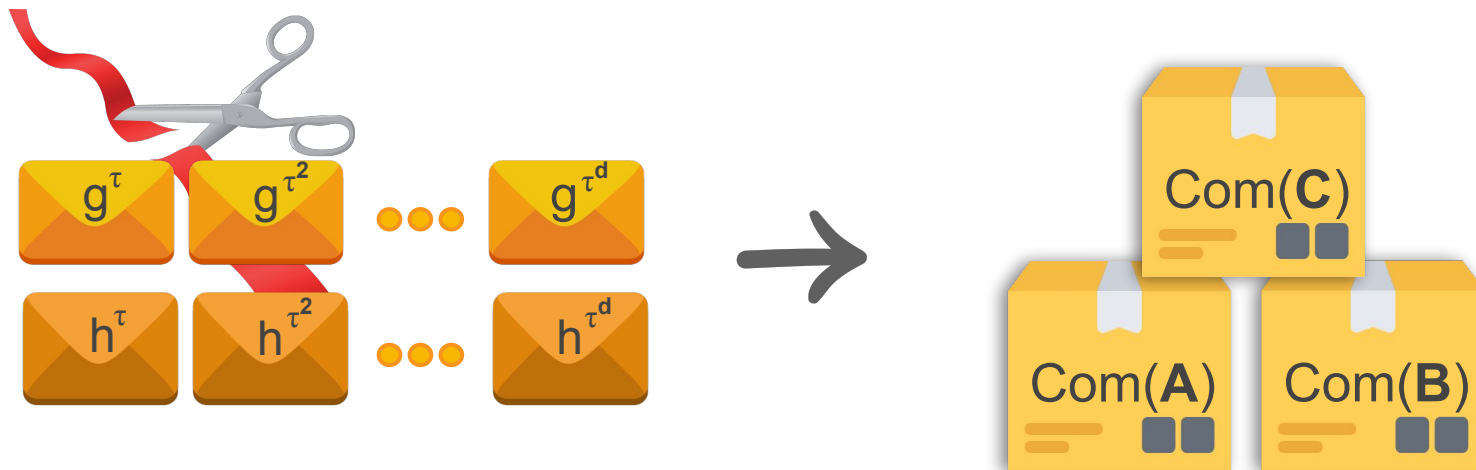
$$A' = A_{\text{left}} A_{\text{right}}$$

A_{final}





Structured Setup





Vector Commitment

KeyGen(λ, n) \rightarrow ck: $[1]_1, [\tau]_1, [\tau^2]_1, \dots, [\tau^{n-1}]_1$

Commit(ck, **a**) \rightarrow  = $[A(\tau)]_1 = A$

$$A(X) = a_1 + a_2 X + a_3 X^2 + \dots + a_n X^{n-1}$$



Vector Commitment

KeyGen(λ, n) \rightarrow ck: $[1]_1, [\tau]_1, [\tau^2]_1, \dots, [\tau^{n-1}]_1$

Commit(ck, **a**) \rightarrow  **a** = $[A(\tau)]_1 = A$

$$A(X) = a_1 + a_2 X + a_3 X^2 + \dots + a_n X^{n-1}$$

Target-Group Commitment

$A =$  A_1  A_2 ...  A_n \rightarrow $Com(A) = e(A_1, [\tau]_2) e(A_2, [\tau^2]_2) \dots e(A_n, [\tau^n]_2)$

Structured Setup

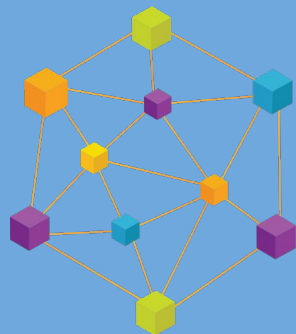


Logarithmic verification for folding commitment key ck



Fast verification ck_{final} :

$$g_{\alpha}(X) = \prod_{i=1}^{\mu} \left(1 + \alpha_{\mu+1-i}^{-2} X^{2^{i-1}} \right)$$

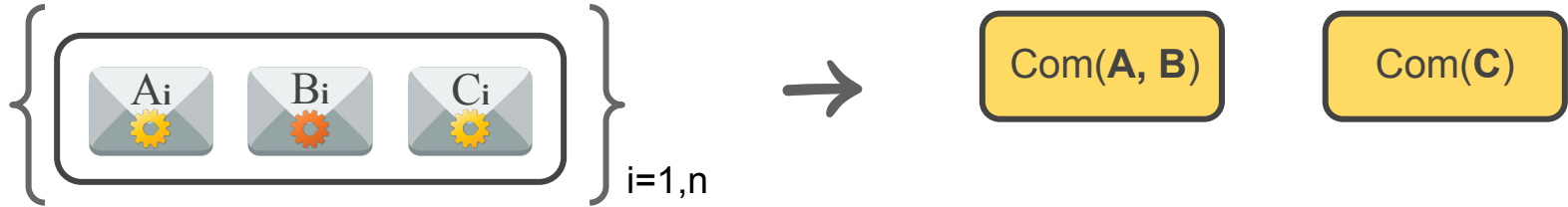


Performance



SNARK Aggregation

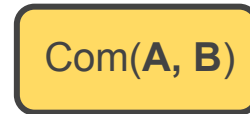
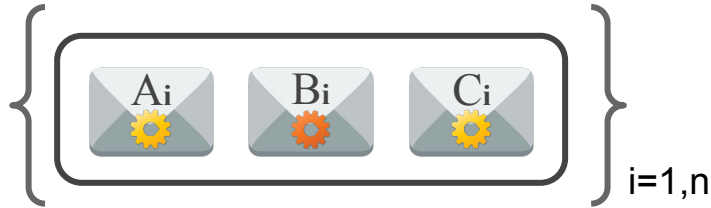
Aggregation





SNARK Aggregation

Aggregation



$$\text{MIPP: } \langle \mathbf{C}, \mathbf{r} \rangle = \prod C_i^{r_i}$$

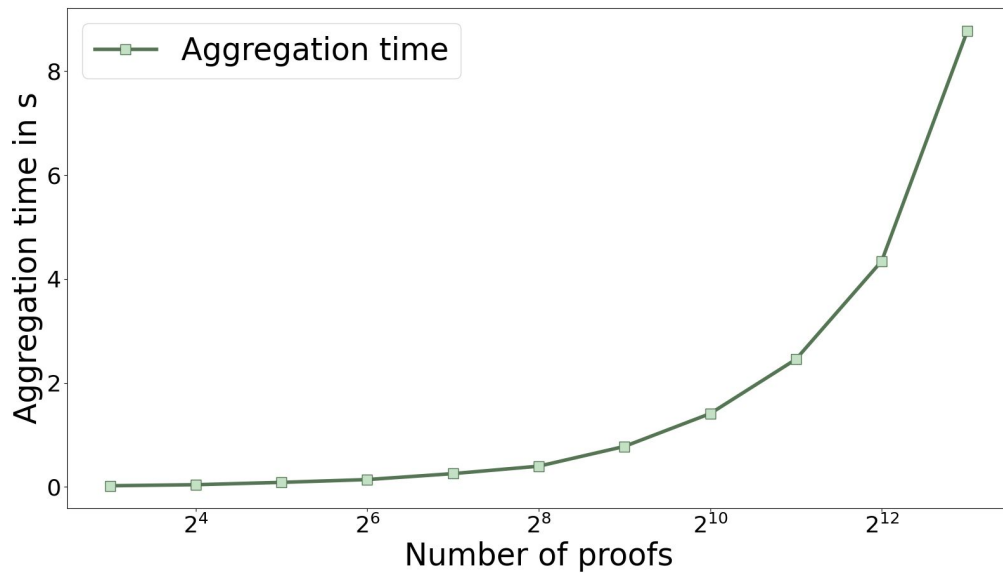


$$\text{TIPP: } \langle \mathbf{A}, \mathbf{B}^{\mathbf{r}} \rangle = \prod e(A_i, B_i^{r_i})$$



log n

Aggregation

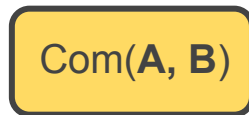
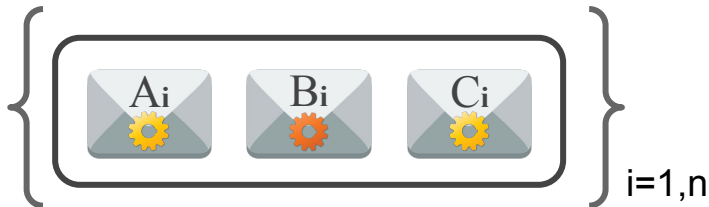


- 8.7s for 8 192 proofs
- Relies heavily on parallelism
- 2^{17} proofs in ~2mn

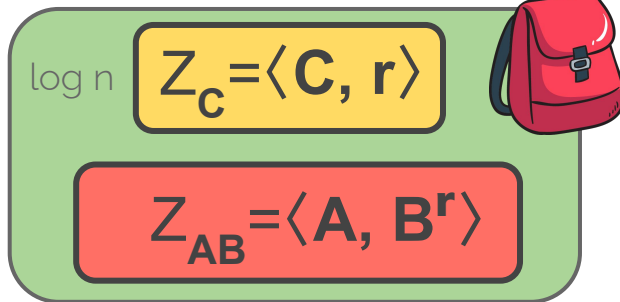


Proof size

Aggregation



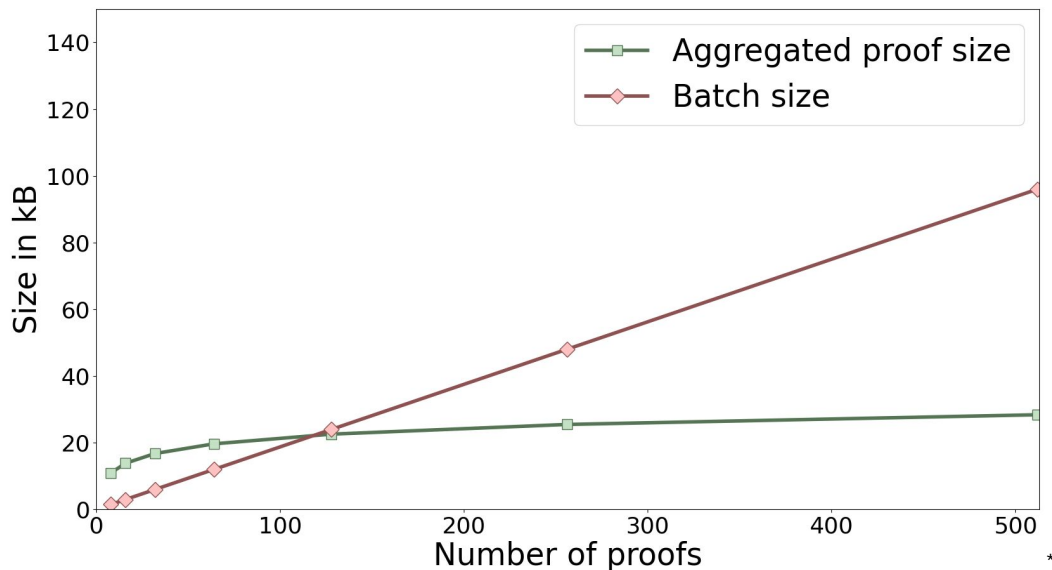
$$\text{MIPP: } \langle \mathbf{C}, \mathbf{r} \rangle = \prod c_i^{r_i}$$



$$\text{TIPP: } \langle \mathbf{A}, \mathbf{B}^{\mathbf{r}} \rangle = \prod e(A_i, B_i^{r_i})$$



Proof size

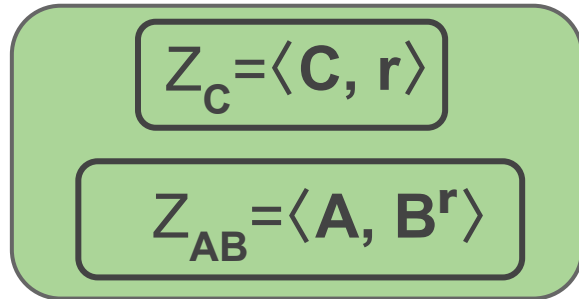
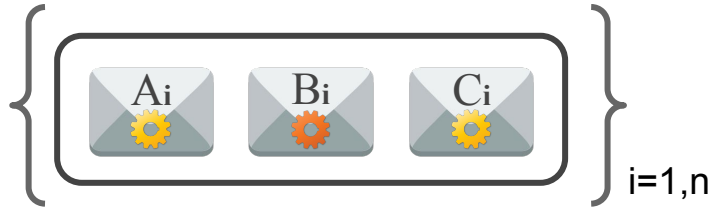


- Use **compression** of target group points
 - based on Torus compression
 - credits - RELIC library implementation
- Turnover at 128 proofs
 - 23kB for aggregated
 - 24kB for "all proofs"

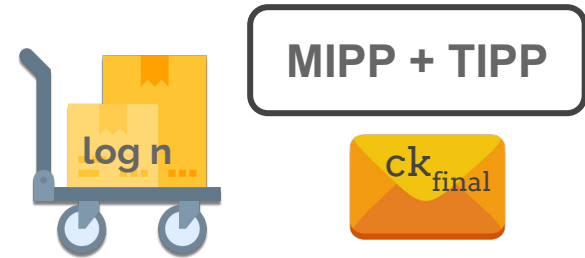


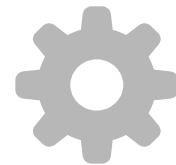
Verification

Aggregation

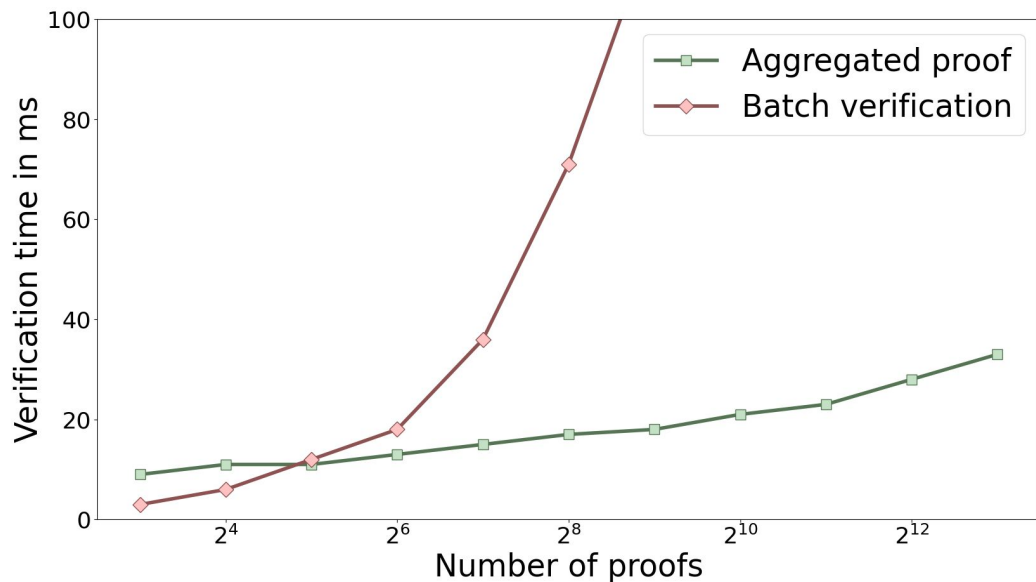


Verification





Verifier Time



- Verifying aggregate proofs becomes **faster** from **32 proofs**
- **8192 proofs** in **33ms**
 - "ratio" of 0.004 ms per proof
- Optimizations:
 - Relies heavily on parallelism
 - MIPP/TIPP combined
 - Batching for pairing checks



SNARK Recursion



Prove many instances



Snark Pack

Proof Size



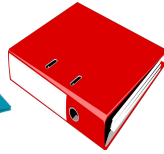
Verifier Time



Prover Time



Recursion

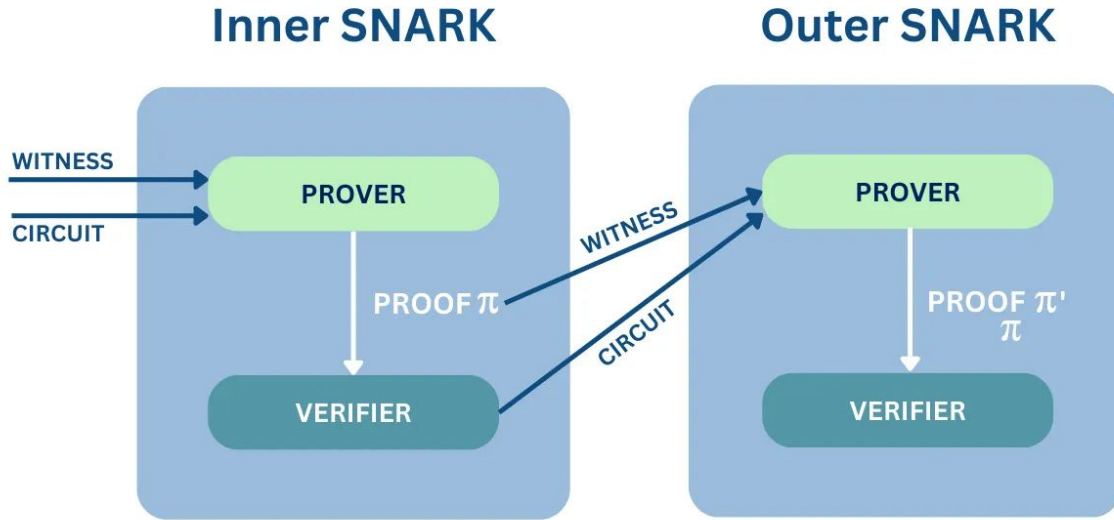


Folding



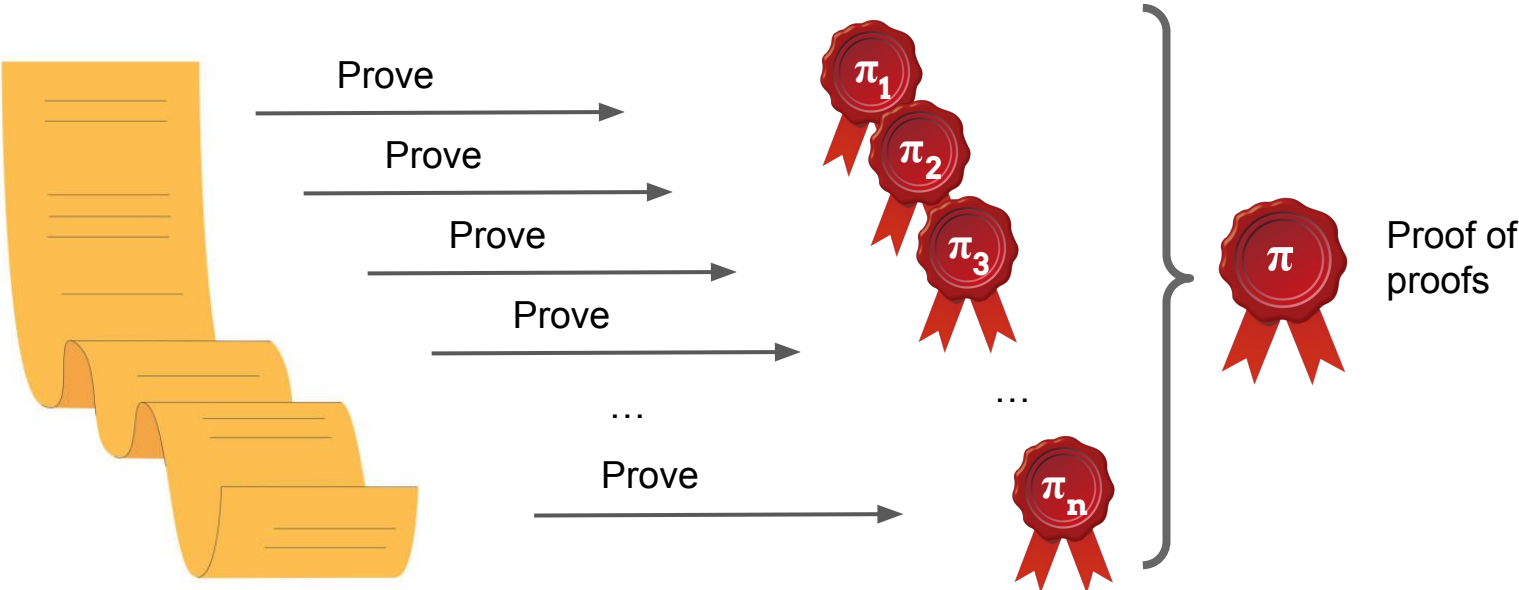


Recursion





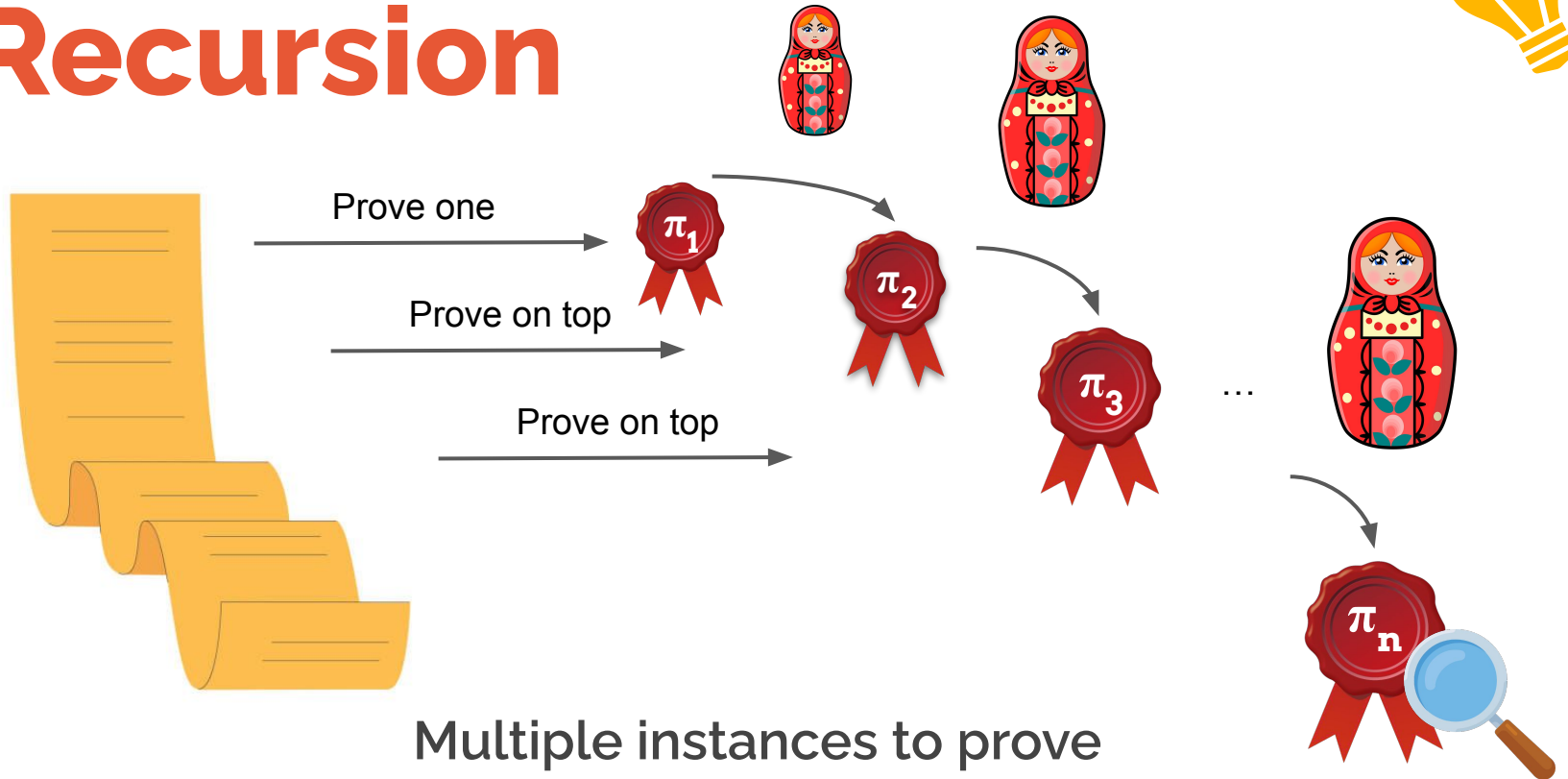
Proof of proofs



Multiple instances to prove



Recursion

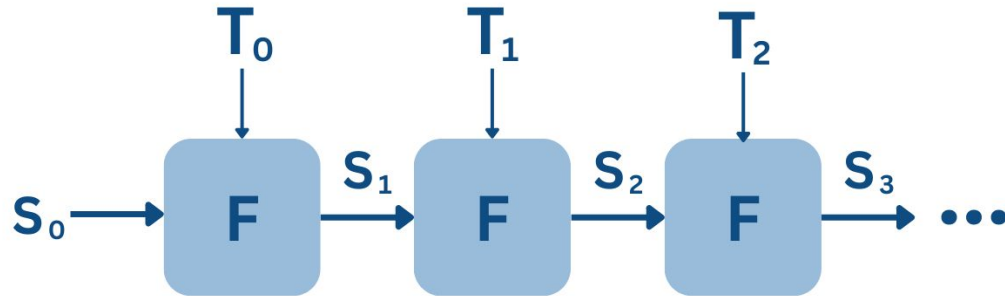


Multiple instances to prove



IVC

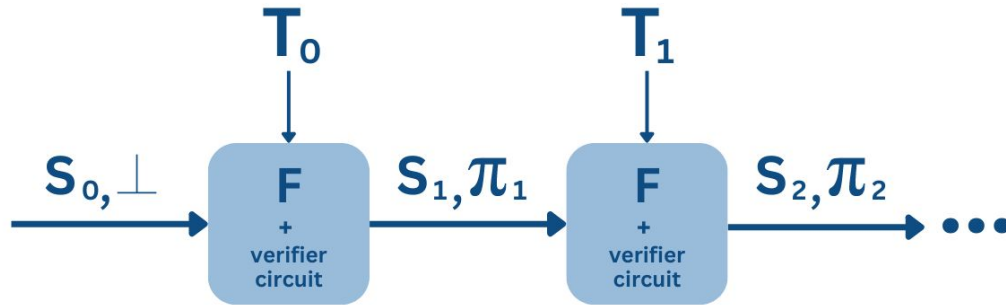
Incremental Verifiable Computation





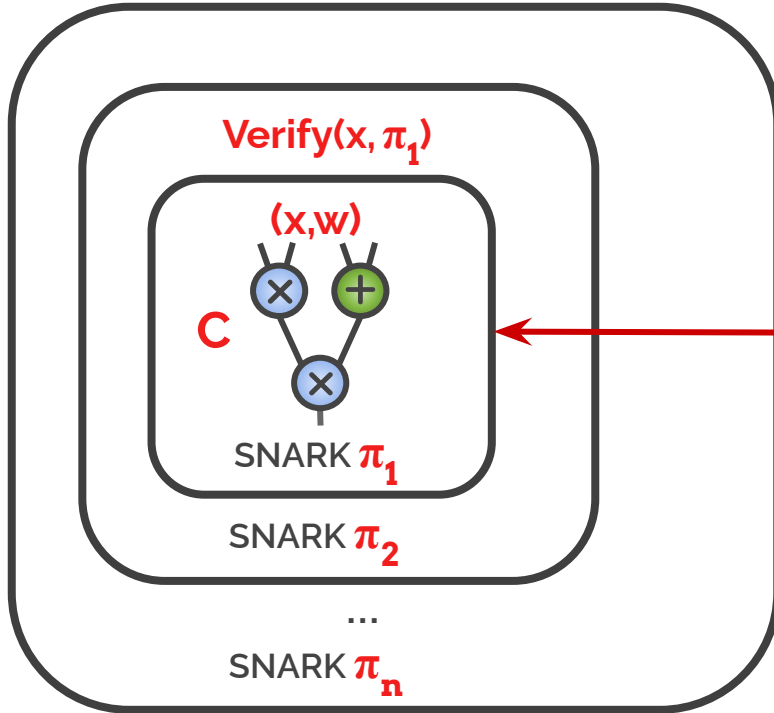
IVC

Incremental Verifiable Computation





Proof Recursion



Verification Runtime

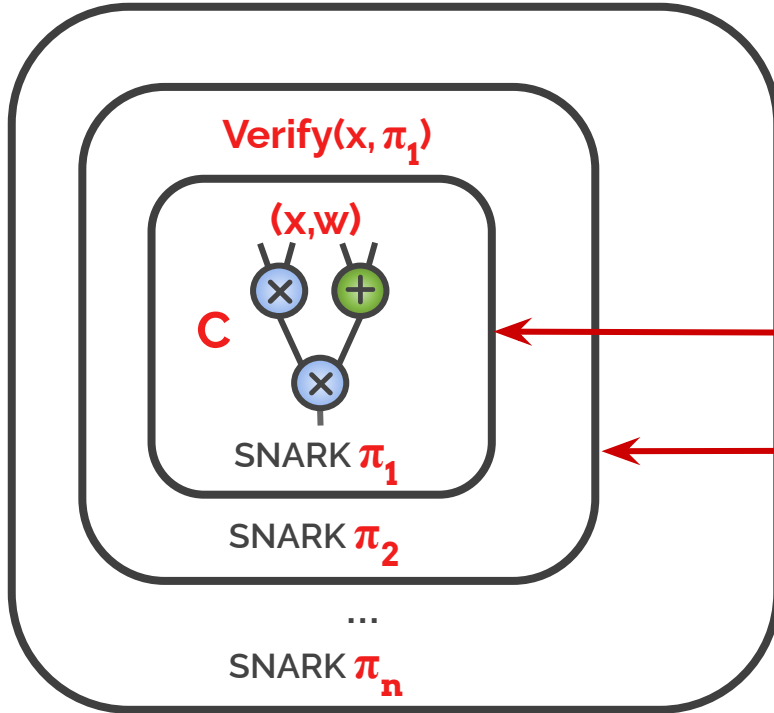
$$|\text{Verify}| = 2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_1}| = 2|C_{\text{SNARK}}|$$

A red arrow points from this equation to the innermost verification box in the diagram.



Proof Recursion



Verification Runtime

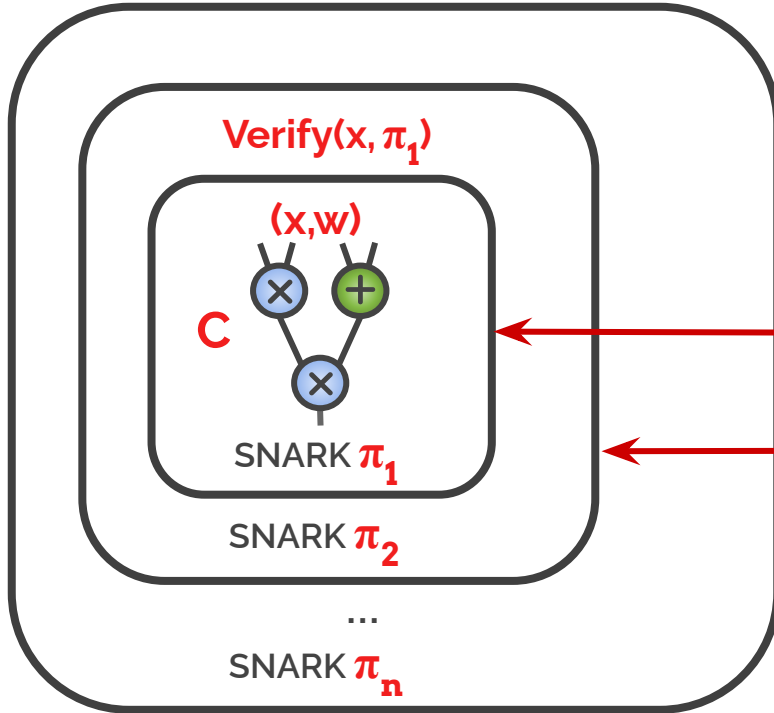
$$|\text{Verify}| = 2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_1}| = 2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_2}| = 2|C_{\text{Verify}_1}| = 2^2|C_{\text{SNARK}}|$$



Proof Recursion



Verification Runtime

$$|\text{Verify}| = 2|C_{\text{SNARK}}|$$

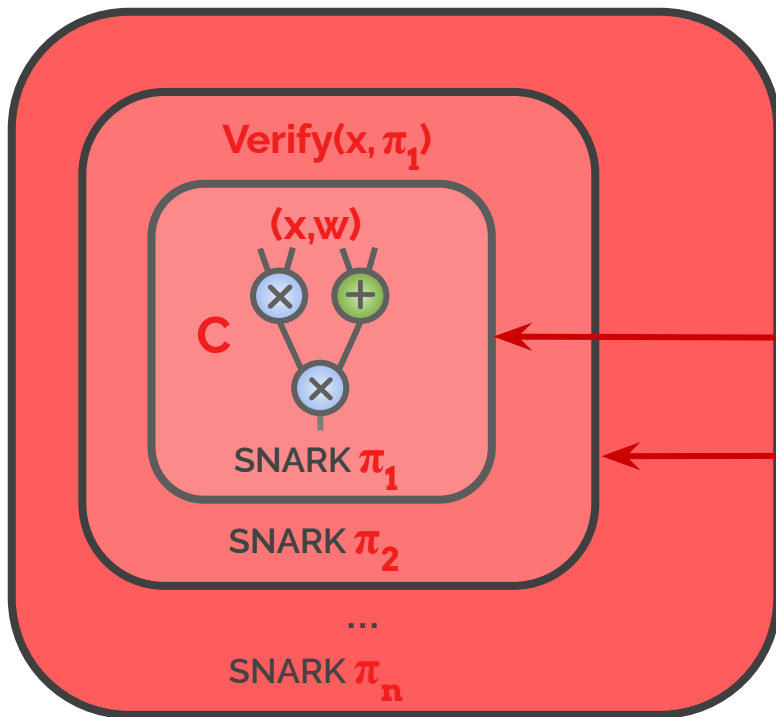
$$|C_{\text{Verify}_1}| = 2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_2}| = 2|C_{\text{Verify}_1}| = 2^2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_3}| = 2|C_{\text{Verify}_2}| = 2^3|C_{\text{SNARK}}|$$



Proof Recursion



Needs **Sublinear** Verification!

$$|C_{\text{Verify}_1}| = 2|C_{\text{SNARK}}|$$

$$|C_{\text{Verify}_2}| = 2|C_{\text{Verify}_1}| = 2^2|C_{\text{SNARK}}|$$

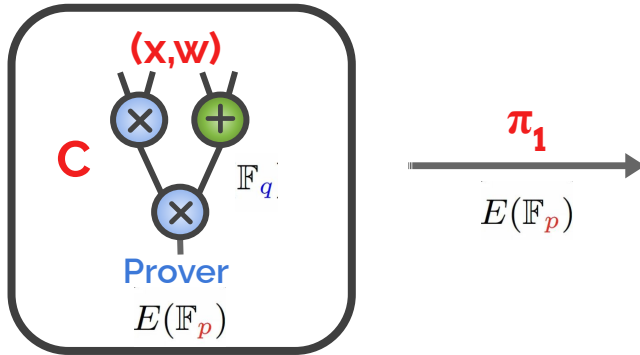
$$|C_{\text{Verify}_3}| = 2|C_{\text{Verify}_2}| = 2^3|C_{\text{SNARK}}|$$



Proof Recursion

Verification Arithmetization

Verify = curve operations

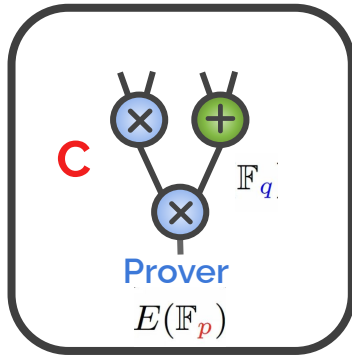




Proof Recursion

Verification Arithmetization

Verify = curve operations



$$\xrightarrow{\pi_1}$$
$$E(\mathbb{F}_p)$$
$$\#E(\mathbb{F}_p) = q$$

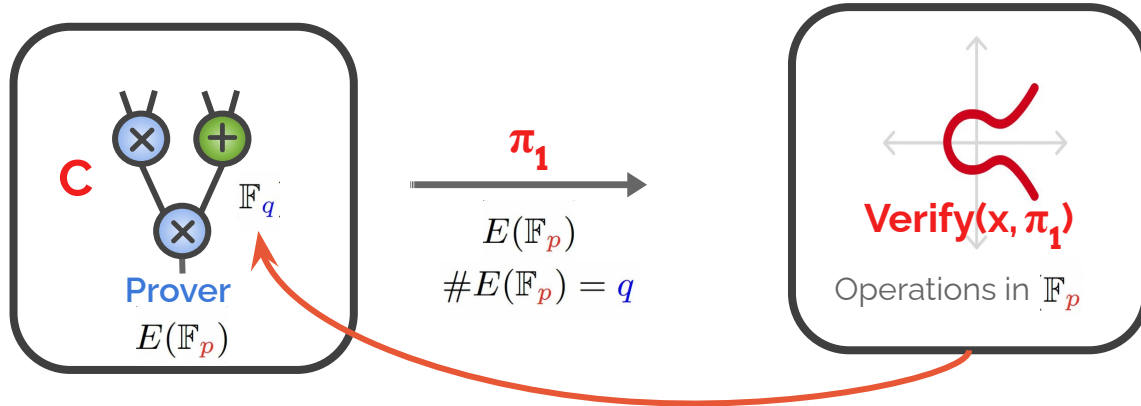




Proof Recursion

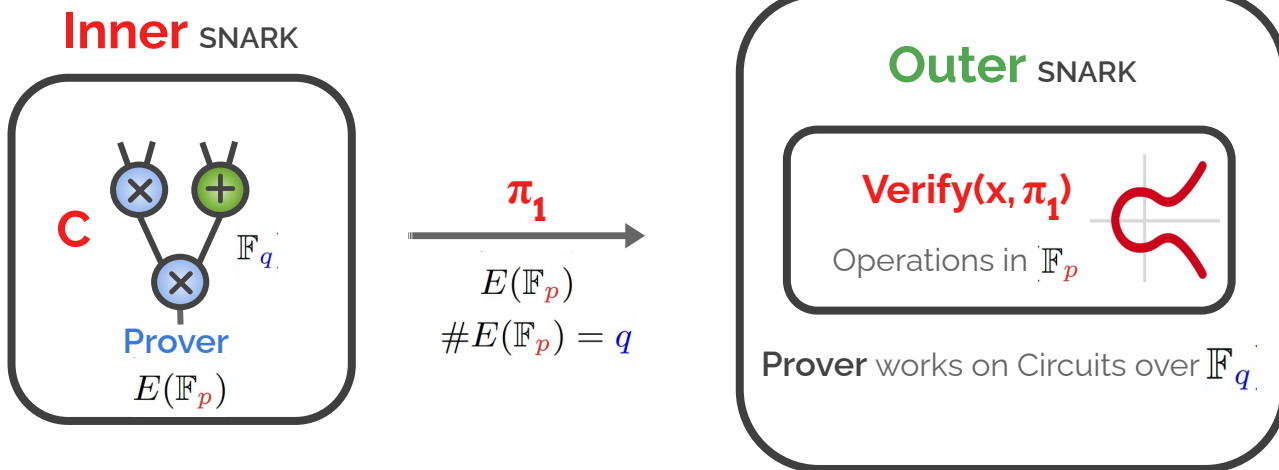
Verification Arithmetization

Verify = curve operations



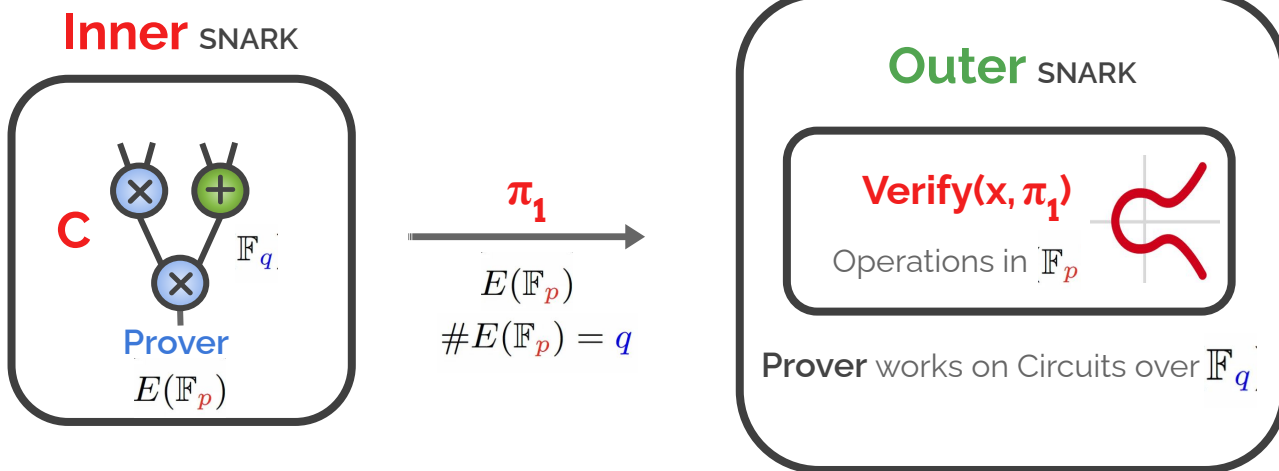


Proof Recursion





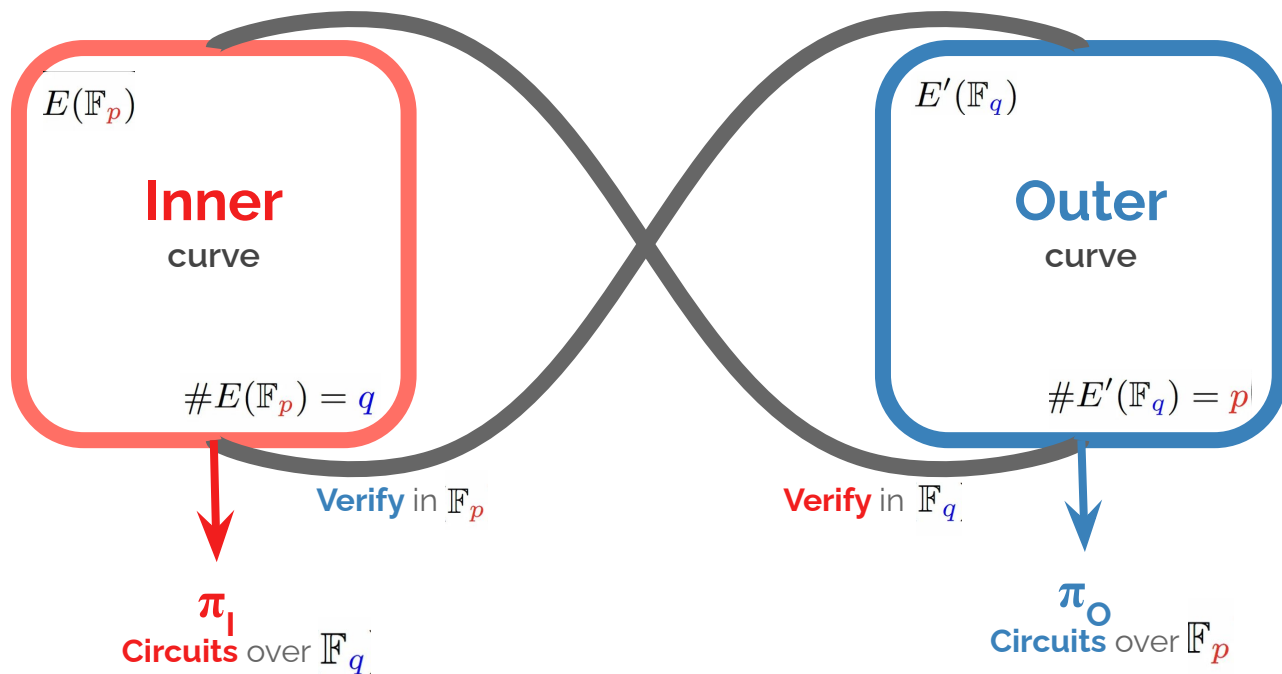
Proof Recursion



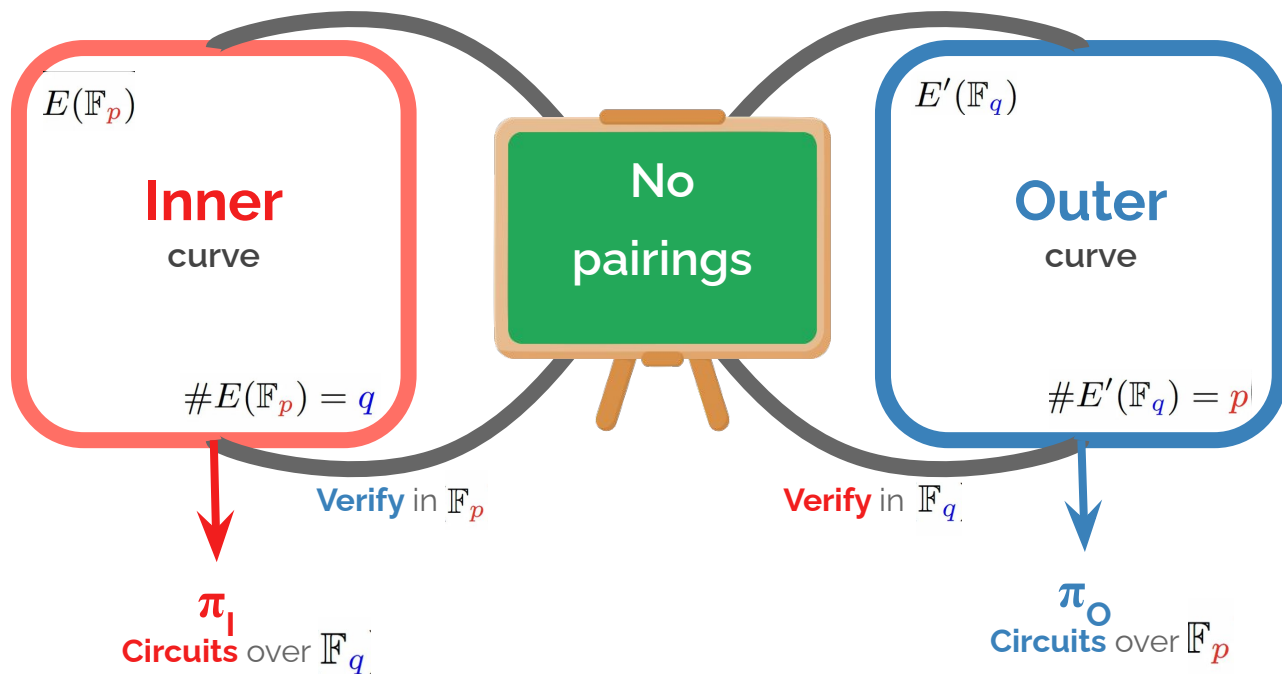
$$p = q$$

Dlog is easy to solve :(

Cycles of Curves

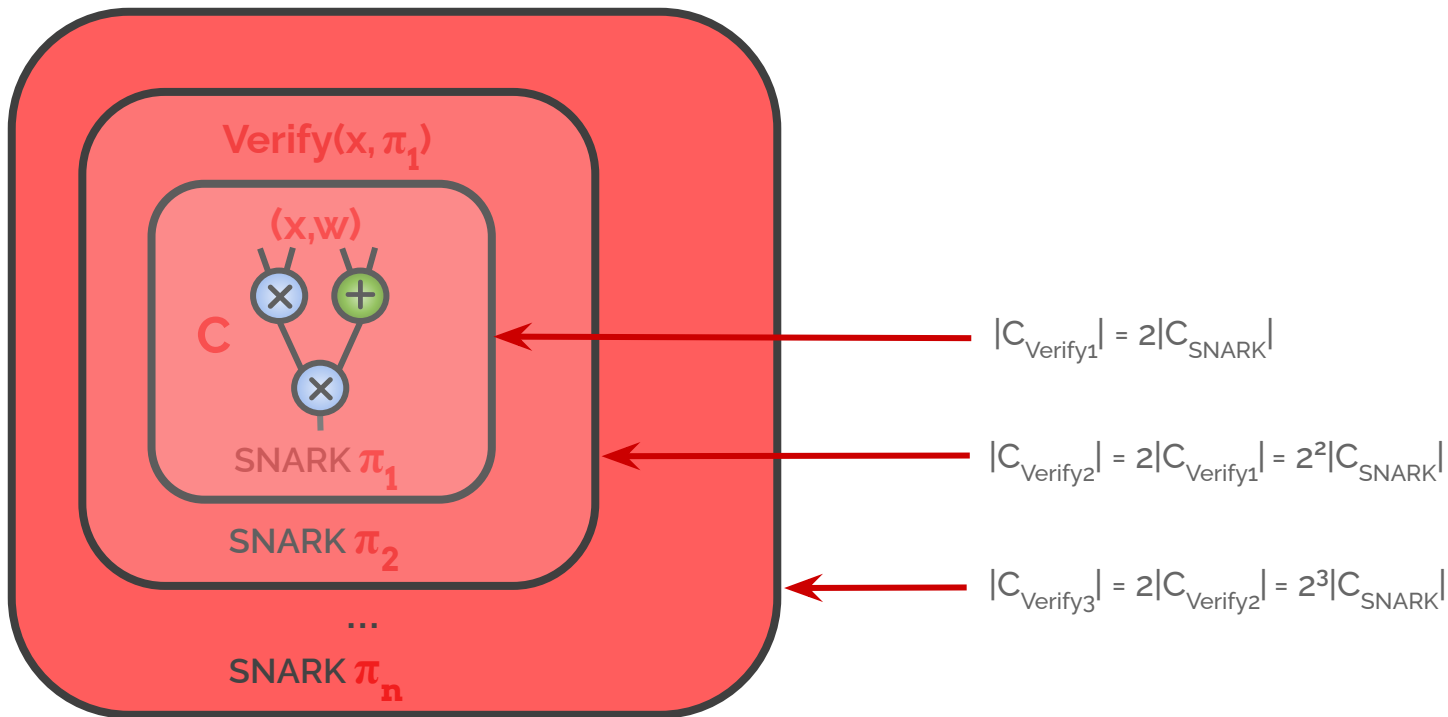


Cycles of Curves



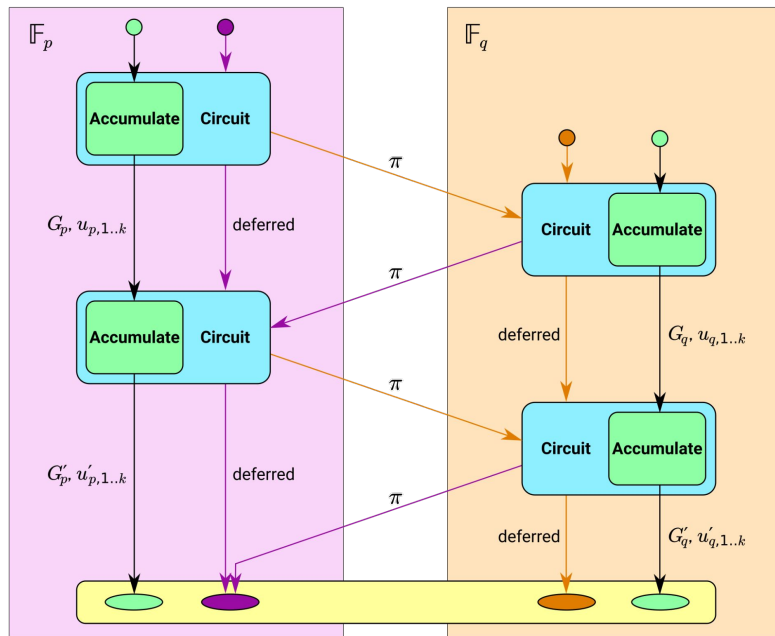


Linear Verification!!!

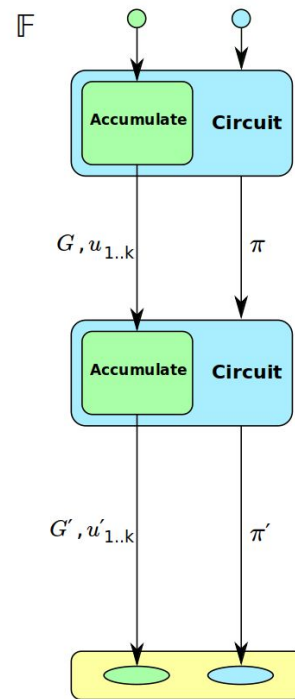




Halo Roadmap



Source: [The Halo2 Book](#)



Thanks



<https://www.di.ens.fr/~nitulesc/>



Credits

Special thanks to all those who made and released these resources for free:

- Presentation template by [SlidesCarnival](#)
- Clip arts by [Iconfinder](#), [Flaticon](#) and [juicy_fish](#)
- Illustrations by [Disneyclips](#)
- Figures from [Blog post](#)